Dual Pivot Quicksort: Verification and Proof using KeY

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Why verify Dual Pivot Quicksort?

Inspired by discovery of Timsort Bug

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Section 1

Algorithm Description





















array index





Why use Dual Pivot Quicksort?

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 Theory: Average number of swaps reduced by 20% (Yaroslavskiy 2009) Why use Dual Pivot Quicksort?

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Why use Dual Pivot Quicksort?

- Theory: Average number of swaps reduced by 20% (Yaroslavskiy 2009)
- Practice: Multi-pivot Quicksorts are more cache-efficient (Kushagra 2014)
- Benchmarking shows it is faster

Java Implementation – Choosing a Sorting Algorithm

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Java Implementation – Single Pivot Partition


Java Implementation – Single Pivot Partition



Java Implementation - Dual Pivot Partition



Java Implementation - Dual Pivot Partition



Java Implementation – Swap Pivot Values Partition



Java Implementation – Swap Pivot Values Partition



















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Section 2

Specification and Proof

Encapsulating source code in its own Java class

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- Subdivision into three classes: One per partitioning style

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- Subdivision into three classes: One per partitioning style
- Writing specification Running KeY Adapting specification or source code

Autopilot Strategy Macro

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- If proof fails:
 - Confirm by generating counterexample
 - Find violated specification condition
 - Adapt specification (or source code)

- Autopilot Strategy Macro
- If proof fails:
 - Confirm by generating counterexample
 - Find violated specification condition
 - Adapt specification (or source code)
- If no proof is found:
 - Increase number of steps (?)
 - Interactive Rule Apps (Quantifier Instantiation, if-then-else-split)
 - Heap Simplification + SMT Solver

- Method extraction
- Exact Localization
- SMT Solver
- Block Contracts

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- Exact Localization
- SMT Solver
- Block Contracts
- Error in specification or lack of resources?

Computation time

- Method extraction
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Localizability

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- Localizability
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- Stability
- Responsiveness

Violation of Single Pivot Partition Invariant

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Violation of Single Pivot Partition Invariant

```
while (a[great] > pivot2) {
    if (great-- == k) {
        break outer;
    }
}
while (a[great] == pivot2) {
    if (great-- == k) {
        break outer;
    }
while (a[great] > pivot) {
    --great;
}
...
```
Violation of Single Pivot Partition Invariant



Section 3

 Verifying a large, complex, real-world Java program with KeY is feasable, but not without challenges

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Statistics – Single Pivot Partition

Method	Nodes	Branches	Time [s]	Rule Apps	Interactive	SMT
case_right	14784	114	17,7	18919	0	0
split	17609	90	23,8	24189	0	0
sort(array, left, right)	18495	101	18,8	22839	0	0
sort(array)	654	7	0,4	1342	0	0
Total	51542	312	60.7	67289	0	0

Statistics – Swap Pivot Values Partition

Method	Nodes	Branches	Time [s]	Rule Apps	Interactive	SMT
move_great_left	1245	16	0,8	2346	0	0
move_less_right	2120	14	1,8	3224	0	0
swap_values	123636	407	246,6	138039	0	0
Total	127001	437	249.2	143609	0	0

Statistics – Dual Pivot Partition

Method	Nodes	Branches	Time [s]	Rule Apps	Interactive	SMT
calc_indices	24533	8	49,6	24835	0	0
insertionsort_indices	50816	365	137,4	73056	0	34
prepare_indices	5332	28	6,4	7153	0	0
move_great_left	1650	15	1,1	2605	0	0
move_great_in_loop	1580	18	1,1	2787	0	0
move_less_right	1928	14	1,4	2967	0	0
loop_body	52134	287	57,3	56263	18	0
split	28751	98	109,6	51666	0	36
sort(int[],left,right)	51342	305	459,6	76973	114	116
sort(int[])	611	5	0,4	1236	0	0
Total	218677	1143	823,9	299541	132	186
Entire Proof	297220	1892	1133,8	510439	132	186