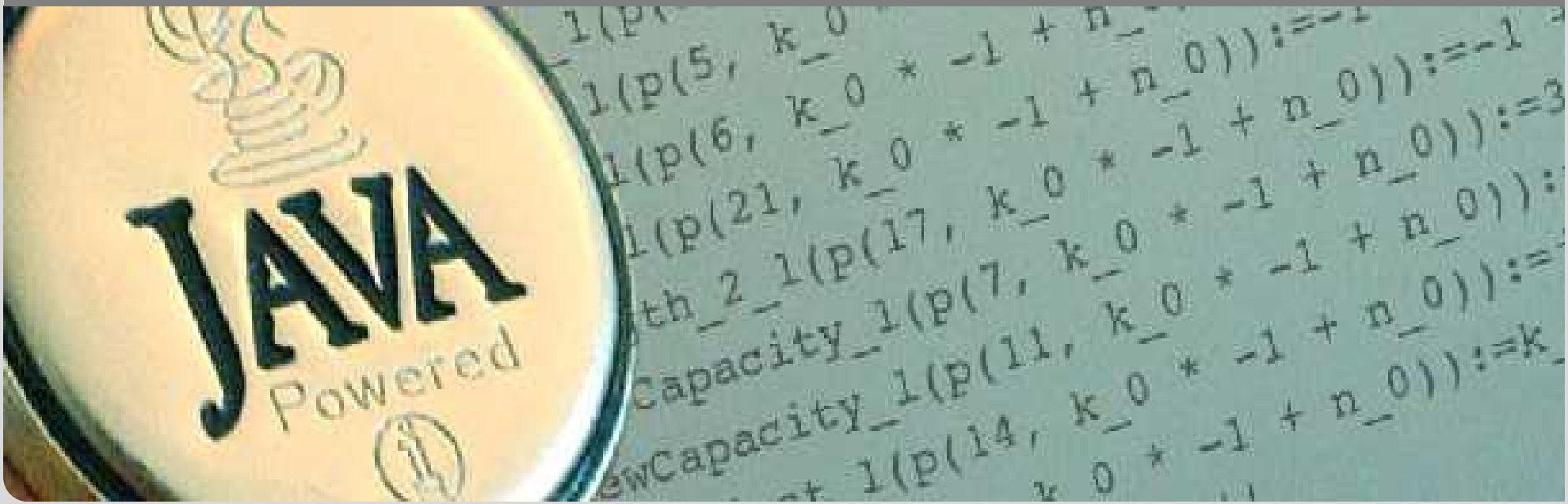
Karlsruhe Institute of Technology

Applications of Formal Verification

Functional Verification of Java Programs: Java Modelling Language

Prof. Dr. Bernhard Beckert · Dr. Vladimir Klebanov | SS 2010

KIT – Institut für Theoretische Informatik



Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

Design by Contract

Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition
then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - *automation* of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,

JML Annotations

```
/*@ public normal_behavior
  @   requires pin == correctPin;
  @   ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

JML Annotations

```
/*@ public normal_behavior
  @   requires pin == correctPin;
  @   ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

JML Annotations

```
/*@ public normal_behavior
  @   requires pin == correctPin;
  @   ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

JML Annotations

```
/*@ public normal_behavior
  @   requires pin == correctPin;
  @   ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

JML Annotations

```
/*@ public normal_behavior
  @   requires pin == correctPin;
  @   ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

JML Annotations

```
/*@ public normal_behavior                    //hello!
  @    requires pin == correctPin;
  @    ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    ...
```

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '*/'.

  $\Rightarrow$ The blue '@'s are not required, but it's a *convention* to use them.
- JML specifications may themselves contain comments

Visibility Modifiers

```
public class ATM {
  private /*@ spec_public @*/ BankCard insertedCard = null;
  private /*@ spec_public @*/
          boolean  customerAuthenticated = false;

  /*@ public normal_behavior ... @*/
```

- Modifiers to specification cases have no influence on their semantics.
- *public* specification items cannot refer to *private* fields.
- Private fields can be declared public for specification purposes only.

Visibility Modifiers

```
public class ATM {
  private /*@ spec_public @*/ BankCard insertedCard = null;
  private /*@ spec_public @*/
          boolean  customerAuthenticated = false;

  /*@ public normal_behavior ... @*/
```

- Modifiers to specification cases have no influence on their semantics.
- *public* specification items cannot refer to *private* fields.
- Private fields can be declared public for specification purposes only.

Visibility Modifiers

```
public class ATM {
  private /*@ spec_public @*/ BankCard insertedCard = null;
  private /*@ spec_public @*/
          boolean  customerAuthenticated = false;

  /*@ public normal_behavior ... @*/
```

- Modifiers to specification cases have no influence on their semantics.
- *public* specification items cannot refer to *private* fields.
- Private fields can be declared public for specification purposes only.

Visibility Modifiers

```
public class ATM {
  private /*@ spec_public @*/ BankCard insertedCard = null;
  private /*@ spec_public @*/
          boolean  customerAuthenticated = false;

  /*@ public normal_behavior ... @*/
```

- Modifiers to specification cases have no influence on their semantics.
- *public* specification items cannot refer to *private* fields.
- Private fields can be declared public for specification purposes only.

Method Contracts

```
/*@ requires r;
  @ assignable a;
  @ diverges d;
  @ ensures post;
  @ signals_only E1,...,En;
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;      //what is the caller's obligation?
  @ assignable a;
  @ diverges d;
  @ ensures post;
  @ signals_only E1,...,En;
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
    normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;      //what is the caller's obligation?
  @ assignable a;    //which locations may be assigned by m?
  @ diverges d;
  @ ensures post;
  @ signals_only E1,...,En;
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
    normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;
  @ signals_only E1,...,En;
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
    normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;  //what must hold on normal termination?
  @ signals_only E1,...,En;
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;  //what must hold on normal termination?
  @ signals_only E1,...,En; //what exc-types may be thrown?
  @ signals(E e) s;
  @*/
T m(...);
```

Abbreviations

```
normal_behavior  =  signals(Exception) false;
exceptional_behavior  =  ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;    //what must hold on normal termination?
  @ signals_only E1,...,En; //what exc-types may be thrown?
  @ signals(E e) s; //what must hold when an E is thrown?
  @*/
T m(...);
```

Abbreviations

```
normal_behavior = signals(Exception) false;
exceptional_behavior = ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;   //what must hold on normal termination?
  @ signals_only E1,...,En; //what exc-types may be thrown?
  @ signals(E e) s; //what must hold when an E is thrown?
  @*/
T m(...);
```

Abbreviations

```
normal_behavior  =  signals(Exception) false;
exceptional_behavior  =  ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;   //what must hold on normal termination?
  @ signals_only E1,...,En; //what exc-types may be thrown?
  @ signals(E e) s; //what must hold when an E is thrown?
  @*/
T m(...);
```

Abbreviations

```
      normal_behavior  =  signals(Exception) false;
exceptional_behavior   =  ensures false;
```

Method Contracts

```
/*@ requires r;       //what is the caller's obligation?
  @ assignable a;     //which locations may be assigned by m?
  @ diverges d;       //when may m non-terminate?
  @ ensures post;   //what must hold on normal termination?
  @ signals_only E1,...,En; //what exc-types may be thrown?
  @ signals(E e) s; //what must hold when an E is thrown?
  @*/
T m(...);
```

Abbreviations

```
      normal_behavior  =  signals(Exception) false;
exceptional_behavior   =  ensures false;
```

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- `instance` invariants *can*, `static` invariants *cannot* refer to `this`
- default: `instance` within classes, `static` within interfaces

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- `instance` invariants *can*, `static` invariants *cannot* refer to `this`
- default: `instance` within classes, `static` within interfaces

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- `instance` invariants *can*, `static` invariants *cannot* refer to `this`
- default: `instance` within classes, `static` within interfaces

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- `instance` invariants *can*, `static` invariants *cannot* refer to `this`
- default: `instance` within classes, `static` within interfaces

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- **`instance`** invariants *can*, **`static`** invariants *cannot* refer to **`this`**
- default: **`instance`** within classes, **`static`** within interfaces

Class Invariants

```
//@ invariant i;
```

- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold “always”
  (cf. *visible state semantics*, *observed state semantics*)
- **`instance`** invariants *can*, **`static`** invariants *cannot* refer to **`this`**
- default: **`instance`** within classes, **`static`** within interfaces

Pure Methods

Pure methods terminate and have no side effects.

After declaring

```
public /*@ pure @*/ boolean cardIsInserted() {
  return insertedCard!=null;
}
```

```
cardIsInserted()
```

could replace

```
insertedCard != null
```

in JML annotations.

Pure Methods

Pure methods terminate and have no side effects.

After declaring

```
public /*@ pure @*/ boolean cardIsInserted() {
  return insertedCard!=null;
}
```

`cardIsInserted()`

could replace

`insertedCard != null`

in JML annotations.

Pure Methods

Pure methods terminate and have no side effects.

After declaring

```
public /*@ pure @*/ boolean cardIsInserted() {
  return insertedCard!=null;
}
```

```
cardIsInserted()
```

could replace

```
insertedCard != null
```

in JML annotations.

Pure Methods

‘`pure`’ $\approx$ ‘`diverges false;`’ + ‘`assignable \nothing;`’

Expressions

- All Java expressions without side-effects
- `==>`, `<==>`: implication, equivalence
- `\forall`, `\exists`
- `\num_of`, `\sum`, `\product`, `\min`, `\max`
- `\old`(...): referring to pre-state in postconditions
- `\result`: referring to return value in postconditions

Expressions

- All Java expressions without side-effects
- ==>, <==>: implication, equivalence
- `\forall`, `\exists`
- `\num_of`, `\sum`, `\product`, `\min`, `\max`
- `\old`(...): referring to pre-state in postconditions
- `\result`: referring to return value in postconditions

Expressions

- All Java expressions without side-effects
- ==>, <==>: implication, equivalence
- `\forall`, `\exists`
- `\num_of`, `\sum`, `\product`, `\min`, `\max`
- `\old`(...): referring to pre-state in postconditions
- `\result`: referring to return value in postconditions

Expressions

- All Java expressions without side-effects
- ==>, <==>: implication, equivalence
- **\forall**, **\exists**
- **\num_of**, **\sum**, **\product**, **\min**, **\max**
- **\old**(...): referring to pre-state in postconditions
- **\result**: referring to return value in postconditions

Expressions

- All Java expressions without side-effects
- ==>, <==>: implication, equivalence
- \forall, \exists
- \num_of, \sum, \product, \min, \max
- \old(...): referring to pre-state in postconditions
- \result: referring to return value in postconditions

Expressions

- All Java expressions without side-effects
- ==>, <==>: implication, equivalence
- **forall**, **exists**
- **num_of**, **sum**, **product**, **min**, **max**
- **old**(...): referring to pre-state in postconditions
- **result**: referring to return value in postconditions

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```

equivalent to

```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```

equivalent to

```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```

equivalent to

```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```

equivalent to

```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Quantification in JML

```
(\forall int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
```

```
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
```
equivalent to
```
(\exists int i; 0<=i && i<\result.length && \result[i]>0)
```

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default `true`.
- JML excludes `null` from the range of quantification.

Generalised and Numerical Quantifiers

`(\num_of C c; e)`	$\#\{c \mid [e]\}$,number of elements of class C with property `e`
`(\sum C c; p; t)`	$\sum_{c:[p]} [t]$
`(\product C c; p; t)`	$\prod_{c:[p]} [t]$
`(\min C c; p; t)`	$min_{c:[p]}\{[t]\}$
`(\max C c; p; t)`	$max_{c:[p]}\{[t]\}$

Generalised and Numerical Quantifiers

`(\num_of C c; e)`	$\#\{c \mid [e]\}$,number of elements of class C with property `e`
`(\sum C c; p; t)`	$\sum_{c:[p]} [t]$
`(\product C c; p; t)`	$\prod_{c:[p]} [t]$
`(\min C c; p; t)`	$min_{c:[p]}\{[t]\}$
`(\max C c; p; t)`	$max_{c:[p]}\{[t]\}$

Generalised and Numerical Quantifiers

`(\num_of C c; e)`	$\#\{c \mid [e]\}$,number of elements of class C with property `e`
`(\sum C c; p; t)`	$\sum_{c:[p]} [t]$
`(\product C c; p; t)`	$\prod_{c:[p]} [t]$
`(\min C c; p; t)`	$min_{c:[p]}\{[t]\}$
`(\max C c; p; t)`	$max_{c:[p]}\{[t]\}$

Generalised and Numerical Quantifiers

`(\num_of C c; e)`	$\#\{c \mid [e]\}$, number of elements of class C with property `e`
`(\sum C c; p; t)`	$\sum_{c:[p]} [t]$
`(\product C c; p; t)`	$\prod_{c:[p]} [t]$
`(\min C c; p; t)`	$min_{c:[p]}\{[t]\}$
`(\max C c; p; t)`	$max_{c:[p]}\{[t]\}$

Generalised and Numerical Quantifiers

`(\num_of C c; e)`	$\#\{c \mid [e]\}$, number of elements of class C with property `e`
`(\sum C c; p; t)`	$\sum_{c:[p]} [t]$
`(\product C c; p; t)`	$\prod_{c:[p]} [t]$
`(\min C c; p; t)`	$\underset{c:[p]}{min}\{[t]\}$
`(\max C c; p; t)`	$\underset{c:[p]}{max}\{[t]\}$

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;   //forbidden (not local, not in assignable)
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;   //forbidden (not local, not in assignable)
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;   //forbidden (not local, not in assignable)
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {




}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;
  tmp.i = 27;
  x = y;
  x.i = 27;
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27;
  x = y;
  x.i = 27;
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;
  x.i = 27;
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;   //forbidden (not local, not in assignable)
}
```

The `assignable` Clauses

Comma-separated list of:

- `e.f` (where `f` a field)
- `a[*]`, `a[x..y]` (where `a` an array expression)
- `\nothing`, `\everything` (default)

Example

```
C x, y;
//@ assignable x, x.i;
void m() {
  C tmp = x;  //allowed (local variable)
  tmp.i = 27; //allowed (in assignable clause)
  x = y;      //allowed (in assignable clause)
  x.i = 27;   //forbidden (not local, not in assignable)
}
```

The `diverges` Clause

```
diverges e;
```

with a boolean JML expression `e` specifies that the method may not terminate only when `e` is true in the pre-state.

Examples

```
diverges false;
```

The method must always terminate.

```
diverges true;
```

The method may terminate or not.

```
diverges n == 0;
```

The method must terminate, when called in a state with `n!=0`.

The `diverges` Clause

```
diverges e;
```

with a boolean JML expression `e` specifies that the method may not terminate only when `e` is true in the pre-state.

Examples

```
diverges false;
```

The method must always terminate.

```
diverges true;
```

The method may terminate or not.

```
diverges n == 0;
```

The method must terminate, when called in a state with `n!=0`.

The `diverges` Clause

```
diverges e;
```

with a boolean JML expression `e` specifies that the method may not terminate only when `e` is true in the pre-state.

Examples

```
diverges false;
```

The method must always terminate.

```
diverges true;
```

The method may terminate or not.

```
diverges n == 0;
```

The method must terminate, when called in a state with `n!=0`.

The `diverges` Clause

```
diverges e;
```

with a boolean JML expression `e` specifies that the method may not terminate only when `e` is true in the pre-state.

Examples

```
diverges false;
```

The method must always terminate.

```
diverges true;
```

The method may terminate or not.

```
diverges n == 0;
```

The method must terminate, when called in a state with `n!=0`.

The `signals` Clauses

```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination $\Rightarrow$ `p` must hold (in post-state)
- exception thrown $\Rightarrow$ must be of type `ET1`, ..., or `ETm`
- exception of type `E1` thrown $\Rightarrow$ `s1` must hold (in post-state)

  ...
- exception of type `En` thrown $\Rightarrow$ `sn` must hold (in post-state)

The `signals` Clauses

```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination $\Rightarrow$ `p` must hold (in post-state)
- exception thrown $\Rightarrow$ must be of type `ET1`, ..., or `ETm`
- exception of type `E1` thrown $\Rightarrow$ `s1` must hold (in post-state)

  ...
- exception of type `En` thrown $\Rightarrow$ `sn` must hold (in post-state)

The `signals` Clauses

```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination $\Rightarrow$ `p` must hold (in post-state)
- exception thrown $\Rightarrow$ must be of type `ET1`, ..., or `ETm`
- exception of type `E1` thrown $\Rightarrow$ `s1` must hold (in post-state)

  ...
- exception of type `En` thrown $\Rightarrow$ `sn` must hold (in post-state)

The `signals` Clauses

```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination $\Rightarrow$ `p` must hold (in post-state)
- exception thrown $\Rightarrow$ must be of type `ET1`, ..., or `ETm`
- exception of type `E1` thrown $\Rightarrow$ `s1` must hold (in post-state)

...

- exception of type `En` thrown $\Rightarrow$ `sn` must hold (in post-state)

The `signals` Clauses

```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination $\Rightarrow$ `p` must hold (in post-state)
- exception thrown $\Rightarrow$ must be of type `ET1`, ..., or `ETm`
- exception of type `E1` thrown $\Rightarrow$ `s1` must hold (in post-state)

  ...
- exception of type `En` thrown $\Rightarrow$ `sn` must hold (in post-state)

Model Fields

```
public interface IBonusCard {

  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?

Remember: There are no attributes in interfaces.

More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

/*@ public instance model int bonusPoints; @*/




  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

/*@ public instance model int bonusPoints; @*/

/*@ ensures bonusPoints == \old(bonusPoints)+newBonusPoints;


  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Model Fields

```
public interface IBonusCard {

/*@ public instance model int bonusPoints; @*/

/*@ ensures bonusPoints == \old(bonusPoints)+newBonusPoints;
  @ assignable bonusPoints;
  @*/
  public void addBonus(int newBonusPoints);

}
```

How to add contracts to abstract methods in interfaces?
Remember: There are no attributes in interfaces.
More precisely: Only static final fields.

Implementing Interfaces

```
public interface IBonusCard {
    /*@ public instance model int bonusPoints; @*/

    /*@ ... @*/
    public void addBonus(int newBonusPoints);
```

Implementation

```
public class BankCard implements IBonusCard{
    public int bankCardPoints;
/*@ private represents bonusPoints = bankCardPoints; @*/

    public void addBonus(int newBonusPoints) {
      bankCardPoints+=newBonusPoints; }
}
```

Implementing Interfaces

```
public interface IBonusCard {
    /*@ public instance model int bonusPoints; @*/

    /*@ ... @*/
    public void addBonus(int newBonusPoints);
```

Implementation

```
public class BankCard implements IBonusCard{
      public int bankCardPoints;


      public void addBonus(int newBonusPoints) {
        bankCardPoints+=newBonusPoints; }
}
```

Implementing Interfaces

```
public interface IBonusCard {
    /*@ public instance model int bonusPoints; @*/

    /*@ ... @*/
    public void addBonus(int newBonusPoints);
```

Implementation

```
public class BankCard implements IBonusCard{
      public int bankCardPoints;
/*@   private represents bonusPoints = bankCardPoints; @*/

      public void addBonus(int newBonusPoints) {
        bankCardPoints += newBonusPoints; }
}
```

Other Representations

```
/*@ private represents bonusPoints
              = bankCardPoints; @*/
```

```
/*@ private represents bonusPoints
              = bankCardPoints * 100; @*/
```

```
/*@ represents x \such_that A(x); @*/
```

Other Representations

```
/*@ private represents bonusPoints
              = bankCardPoints; @*/
```

```
/*@ private represents bonusPoints
              = bankCardPoints * 100; @*/
```

```
/*@ represents x \such_that A(x); @*/
```

Other Representations

```
/*@ private represents bonusPoints
              = bankCardPoints; @*/
```

```
/*@ private represents bonusPoints
              = bankCardPoints * 100; @*/
```

```
/*@ represents x \such_that A(x); @*/
```

Inheritance of Specifications in JML

- An invariant to a class is inherited by all its subclasses.
- An operation contract is inherited by all overridden methods.
  It can be extended there.

Inheritance of Specifications in JML

- An invariant to a class is inherited by all its subclasses.
- An operation contract is inherited by all overridden methods.
  It can be extended there.

Inheritance of Specifications in JML

- An invariant to a class is inherited by all its subclasses.
- An operation contract is inherited by all overridden methods.
  It can be extended there.

Other JML Features

- assertions '`//@ assert e;`'
- loop invariants '`//@ maintaining p;`'
- data groups
- **refines**
- many more...

Other JML Features

- assertions '`//@ assert e;`'
- loop invariants '`//@ maintaining p;`'
- data groups
- **refines**
- many more...

Other JML Features

- assertions '`//@ assert e;`'
- loop invariants '`//@ maintaining p;`'
- data groups
- refines
- many more...

Other JML Features

- assertions '`//@ assert e;`'
- loop invariants '`//@ maintaining p;`'
- data groups
- **refines**
- many more...

Other JML Features

- assertions '`//@ assert e;`'
- loop invariants '`//@ maintaining p;`'
- data groups
- **refines**
- many more...

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '`invariant x != null;`'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'`requires p != null;`'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'`ensures \result != null;`'

non_null is the default!
If something may be `null`, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** **result** != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** **result** != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** \result != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** \result != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** \result != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** **result** != **null**;'

non_null is the default!
If something may be null, you have to declare it **nullable**

Nullity

JML has modifiers **non_null** and **nullable**

```
private /*@spec_public non_null@*/ Object x;
```

⇝ implicit invariant added to class: '**invariant** x != **null**;'

```
void m(/*@non_null@*/ Object p);
```

⇝ implicit precondition added to all contracts:
'**requires** p != **null**;'

```
/*@non_null@*/ Object m();
```

⇝ implicit postcondition added to all contracts:
'**ensures** **result** != **null**;'

non_null is the default!
If something may be `null`, you have to declare it **nullable**

Problems with Specifications Using Integers

```
/*@ requires y >= 0;
  @ ensures
  @  \result * \result <= y &&
  @ y < (abs(\result)+1) * (abs(\result)+1);
  @ */
  public static int isqrt(int y)
```

For $y = 1$ and $\backslash result = 1073741821 = \frac{1}{2}(max\_int - 5)$ the above postcondition is true, though we do not want 1073741821 to be a square root of 1.
JML uses the Javasemantics of integers:

$$1073741821 * 1073741821 = -2147483639$$
$$1073741822 * 1073741822 = 4$$

Problems with Specifications Using Integers

```
/*@ requires y >= 0;
  @ ensures
  @  \result * \result <= y &&
  @ y < (abs(\result)+1) * (abs(\result)+1);
  @ */
  public static int isqrt(int y)
```

For $y = 1$ and $\backslash result = 1073741821 = \frac{1}{2}(max\_int - 5)$ the above postcondition is true, though we do not want 1073741821 to be a square root of 1.

JML uses the Javasemantics of integers:

$$1073741821 * 1073741821 = -2147483639$$
$$1073741822 * 1073741822 = 4$$

Problems with Specifications Using Integers

```
/*@ requires y >= 0;
  @ ensures
  @  \result * \result <= y &&
  @ y < (abs(\result)+1) * (abs(\result)+1);
  @ */
  public static int isqrt(int y)
```

For $y = 1$ and $\backslash result = 1073741821 = \frac{1}{2}(max\_int - 5)$ the above postcondition is true, though we do not want 1073741821 to be a square root of 1.
JML uses the Javasemantics of integers:

$$\begin{aligned} 1073741821 * 1073741821 &= -2147483639 \\ 1073741822 * 1073741822 &= 4 \end{aligned}$$

Problems with Specifications Using Integers

```
/*@ requires y >= 0;
  @ ensures
  @  \result * \result <= y &&
  @ y < (abs(\result)+1) * (abs(\result)+1);
  @ */
  public static int isqrt(int y)
```

For $y = 1$ and $\backslash result = 1073741821 = \frac{1}{2}(max\_int - 5)$ the above postcondition is true, though we do not want 1073741821 to be a square root of 1.
JML uses the Javasemantics of integers:

$$\begin{aligned} 1073741821 * 1073741821 &= -2147483639 \\ 1073741822 * 1073741822 &= 4 \end{aligned}$$

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing

JML Tools

Many tools support JML (see JML homepage). Among them:

- `jml`: JML syntax checker
- `jmldoc`: code documentation (like Javadoc)
- `jmlc`: compiles Java+JML into bytecode with assertion checks
- `jmlunit`: unit testing (like JUnit)
- `rac`: runtime assertion checker
- ESC/Java2: leightweight static verification
- KeY: full static verification
- `OpenJML`: tool suite, under development

The tools do not yet support the new features of Java 5!
e.g.: no generics, no enums, no enhanced for-loops, no autoboxing