

Formale Systeme II: Theorie Axiomatic Set Theory

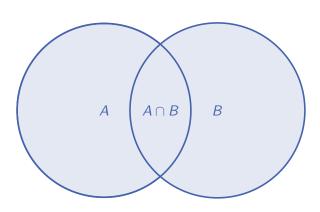
SS 2016

Prof. Dr. Bernhard Beckert · Dr. Mattias Ulbrich Slides partially courtesy by Prof. Dr. Peter H. Schmitt

Motivation

Do you know set theory?





Do you know axiomatic set theory?



$$\forall z(z \in x \leftrightarrow z \in y) \rightarrow x = y.$$

$$\exists y(y \in x) \rightarrow \exists y(y \in x \land \forall z \neg (z \in x \land z \in y)).$$

$$\exists y \forall z(z \in y \leftrightarrow z \in x \land \phi(z)).$$
for any formula ϕ not containing y .
$$\exists y \forall x(x \notin y).$$

$$\exists y \forall x(x \in y \leftrightarrow x = z_1 \lor x = z_2).$$

$$\exists y \forall z(z \in y \leftrightarrow \forall u(u \in z \rightarrow u \in x)).$$

$$\exists y \forall z(z \in y \leftrightarrow \exists u(z \in u \land u \in x)).$$

$$\exists w(\emptyset \in w \land \forall x(x \in w \rightarrow \exists z(z \in w \land \psi \cup x)).$$

$$\forall x, y, z(\psi(x, y) \land \psi(x, z) \rightarrow y = z) \rightarrow \exists u \forall w_1(w_1 \in u \leftrightarrow \exists w_2(w_2 \in a \land \psi(w_2, w_1))).$$

$$\forall x(x \in z \rightarrow x \neq \emptyset \land \forall y(y \in z \rightarrow x \cap y = \emptyset \lor x = y))$$

$$\rightarrow \exists u \forall x \exists v(x \in z \rightarrow u \cap x = \{v\}).$$

Georg Cantor





Georg F.L.P. Cantor

1918

1845 born in St. Petersburg
1862 studies in Zürich, Göttingen
1867 and Berlin
1872 foundations of
1884 axiomatic set theory
1869 Professor
1918 in Halle (Saale)

died in Halle

Georg Cantor



Pioneering Publications:

Über unendliche Punctmanichfaltigkeiten.

Math. Ann. 15(1879), 1–7, 17(1880), 355–358, 20(1882), 113–121, 21(1883), 51–58 and 545–586, 23(1884), 453–488

Beiträge zur Begründung der transfiniten Mengenlehre.

Math. Ann. 46(1895), 481–512, 49(1897), 207–245.

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'Naive' Set Theory



Beiträge zur Begründung der transfiniten Mengenlehre.

Von

GEORG CANTOR in Halle a./S.

(Erster Artikel.)

"Hypotheses non fingo."

"Neque enim leges intellectui aut rebus damus ad arbitrium nostrum, sed tanquam scribae fideles ab ipsius naturae voce latas et prolatas excipimus et describimus."

"Veniet tempus, quo ista quae nunc latent, in lucem dies extrahat et longioris aevi diligentia."

§ 1.



Der Mächtigkeitsbegriff oder die Cardinalzahl.

Unter einer "Menge" verstehen wir jede Zusammenfassung M von bestimmten wohlunterschiedenen Objecten m unsrer Anschauung oder unseres Denkens (welche die "Elemente" von M genannt werden) zu einem Ganzen.

In Zeichen drücken wir dies so aus:

 $(1) M = \{m\}.$



Naive Set Theory is not consistent



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Cantor's Antinomy

The "set of all conceivable objects" cannot exist: Its powerset would have to have larger cardinality.



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Russell's Antinomy (1903)

Let $R := \{x \mid x \notin x\}$. Now $R \in R$ is neither true nor false.

⇒ Sudden end for naive set theory.



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Let $R := \{x \mid x \notin x\}$. Now $R \in R$ is neither true nor false. \implies Sudden end for naive set theory.

Insight:

A class term $\{x \mid \varphi(x)\}$ does not necessarily describe a set!

(Historical Sidenote)



Gottlieb Frege (1903): Grundgesetze der Arithmetik, Nachwort

Einem wissenschaftlichen Schriftsteller kann kaum etwas Unerwünschteres begegnen, als daß ihm nach Vollendung einer Arbeit eine der Grundlagen seines Baues erschüttert wird. In diese Lage wurde ich durch einen Brief des Herrn Bertrand Russell versetzt, als der Druck dieses Bandes sich seinem Ende näherte.



Logicism (G. Frege)



Philosophical views:

What is the relationship betwen logics and mathematics?

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Philosophical views:

What is the relationship betwen **logics** and **mathematics**?

Logical reasoning is a branch of mathematics.
 Mathematical subjects are "there" and wait to be described, formally captured.

Logicism (G. Frege)



Philosophical views:

What is the relationship betwen logics and mathematics?

Logical reasoning is a branch of mathematics.
 Mathematical subjects are "there" and wait to be described, formally captured.

or

- Mathematics is an application of logical reasoning.
 - There are valid axioms which are evidently true.
 - All true propositions must be formally derived from axioms.

This lecture



- History of set theory, logicism
- Zermelo-Fraenkel as a prominent first order theory
- Zermelo-Fraenkel as an example of modelling in FOL
- Zermelo-Fraenkel as foundations of mathematics

First Order Logic – Conservative Extension

Conservative Extension



Definition (proof-theoretic)

Let $\Sigma_1 \subseteq \Sigma_2$ be signatures, and T_i set of sentences in Fml_{Σ_i} . T_2 is called a **conservative extension** of T_1 if

$$T_1 \models \varphi \iff T_2 \models \varphi$$
 for all sentences $\varphi \in \mathit{Fml}_{\Sigma_1}$

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Sufficient criterion (model-theoretic)

• Every model for T_1 can be extended to a model of T_2 .

and

• Every restriction of a model of T_2 is a model of T_1 .



Let
$$\Sigma_0 = \{(0, s), (=), \alpha\}$$

Axioms T0

- $\forall x. \ \neg s(x) = 0$
- $\forall x, y. \ s(x) = s(y) \rightarrow x = y$



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Axioms T2:
$$\Sigma_2 = \Sigma_1$$

- Axioms T2
- $x = 0 \lor \exists y.x = s(y)$



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- Axioms T0
 - 1 = s(0)

Axioms T2:
$$\Sigma_2 = \Sigma_1$$

not conservative extension of T1

- Axioms T2
- $x = 0 \lor \exists y.x = s(y)$

Conservativity Theorem



Theorem

Let Σ be a signature, T a Σ -theory and $\varphi(x, \bar{y})$ a Σ -formula. Let $f \notin \Sigma$ be new function symbol

If
$$T \models \forall \bar{y}. \exists x. \varphi(x, \bar{y})$$

then

$$T \cup \{ \forall y. \ \varphi(f(\bar{y}), y) \text{ is a conservative extension of } T \text{ (over } \Sigma \cup \{f\}) \}$$

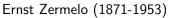
Example

If $\forall y \exists x. \ y = x \cdot x$ is a theorem of some theory R+, then function symbol sqrt can be added as conservative extension to R+ with definition $\forall y. \ y = sqrt(y) \cdot sqrt(y)$.

Zermelo and Fraenkel









Abraham Fraenkel (1891-1965)

Zermelo and Fraenkel







Ernst Zermelo (1871-1953) Abraham Fraenkel (1891-1965)

1907 Zermelo proposes an axiom system with 7 axioms

1921 Fraenkel adds the replacement axiom

1930 Zermolo adds the foundation axiom Axiom of choice was in initial set

Signature



$$\Sigma = \{F, P, \alpha\}$$
 with

- $F = \emptyset$
- $P = \{ \in, = \}$
- $\alpha(\in) = \alpha(=) = 2$

The semantics of equality is "built in", as usual.

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That's it. ...

Only two predicate symbols in the signature.

All other symbols often used $(\emptyset, \cup, \subset, \ldots)$ are derived symbols.

ZF



We will look at the axioms individually:

- Original textual formulation, from Ernst Zermelo: Untersuchungen über die Grundlagen der Mengenlehre. In: Mathematische Annalen. 65 (1908)
- As FOL formulas over the above signature, in modern notation



A1: Extensionality

,,Ist jedes Element einer Menge M gleichzeitig Element der Menge N und umgekehrt $[\ldots]$, so ist immer M=N.

Oder kürzer: jede Menge ist durch ihre Elemente bestimmt." [Zermelo, 1907]

$$\forall z. \ (z \in x \leftrightarrow z \in y) \to x = y$$



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- Free variables in axioms are implicitly universally quantified.
- What about the converse implication?
 (Hint: Remember semantics of "="!)



A2: Foundation / Regularity

"Jede (rückschreitende) Kette von Elementen, in welcher jedes Glied Element des vorangehenden ist, bricht mit endlichem Index ab […].

Oder, was gleichbedeutend ist: Jeder Teilbereich T enthält wenigstens ein Element t_0 , das kein Element t in T hat." [Zermelo, 1930]

$$(\exists y.\ y \in x) \rightarrow \exists y.\ (y \in x \land \forall z.\ \neg(z \in x \land z \in y))$$



A3: Separation Schema

"Ist die Klassenaussage F(x) definit* für alle Elemente einer Menge M, so besitzt M immer eine Untermenge M_F , welche alle diejenigen Elemente x von M, für welche F(x) wahr ist, und nur solche als Elemente enthält."

 $* \approx F(x)$ ist eine Formel.

[Zermelo, 1908]

$$\exists y. \forall z. (z \in y \leftrightarrow z \in x \land \phi(z))$$

for any formula ϕ not containing y.

• is an axiom *schema*, contains a placeholder symbol



A4: Empty set

,,Es gibt eine (uneigentliche) Menge, die Nullmenge O, welche gar keine Elemente enthält." [Zermelo, 1908]

$$\exists y. \forall x. \neg (x \in y).$$



A5: Pair set

,,[...]; sind a, b irgend zwei Dinge des Bereiches, so existiert immer eine Menge $\{a,b\}$, welche sowohl a als b, aber kein von beiden verschiedenes Ding x als Element enthält." [Zermelo, 1908]

$$\exists y. \forall x. (x \in y \leftrightarrow x = z_1 \lor x = z_2)$$



A6: Power set

,,Jeder Menge T entspricht eine zweite Menge UT (die Potenzmenge von T), welche alle Untermengen yon T und nur solche als Elemente enthält." [Zermelo, 1908]

$$\exists y. \forall z. (z \in y \leftrightarrow \forall u (u \in z \rightarrow u \in x))$$



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- Countable infinite set ⇒ uncountable power set expected.
- Löwenheim-Skolem: There is a countable model of set theory.
- ⇒ Not all subsets can be guaranteed to exist



A7: Union / Sum

Jeder Menge T entspricht eine Menge GT (die Vereinigungsmenge von T), welche alle Elemente der Elemente yon T und nur solche als Elemente enthält.

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$$\exists y \forall z (z \in y \leftrightarrow \exists u (z \in u \land u \in x))$$

$$GT = \bigcup T$$



A8: Infinity

Different Notion

"Der Bereich* enthält mindestens eine Menge Z, welche die Nullmenge als Element enthält und so beschaffen ist, daß jedem ihrer Elemente a ein weiteres Element der Form $\{a\}$ entspricht, oder welche mit jedem ihrer Elemente a auch die entsprechende Menge $\{a\}$ als Element enthält."

* Universum/Domäne

[Zermelo, 1907]

$$\exists w. (\emptyset \in w \land \forall x (x \in w \rightarrow \exists z (z \in w \land \forall u (u \in z \leftrightarrow u \in x \lor u = x))))$$



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 $\emptyset \in Z \land \forall a. (a \in Z \rightarrow \{a\} \in Z)$



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- $\emptyset \in Z \land \forall a. (a \in Z \rightarrow \{a\} \in Z)$
- $\emptyset \in Z \land \forall a. (a \in Z \rightarrow a \cup \{a\} \in Z)$



A9: Replacement

"Ist M eine Menge und wird jedes Element von M durch ein "Ding des Bereiches" […] ersetzt, so geht M wiederum in eine Menge über." [Fraenkel, 1921]

$$\forall x, y, z. \ (\psi(x, y) \land \psi(x, z) \rightarrow y = z) \rightarrow \\ \exists u. \forall w_1. \ (w_1 \in u \leftrightarrow \exists w_2 (w_2 \in a \land \psi(w_2, w_1)))$$



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$$(\forall x. \ x \in u \to \exists y. \psi(x, y)) \land (\forall x, y, z. \ (\psi(x, y) \land \psi(x, z) \to y = z)) \to \exists v. \forall y. (y \in v \leftrightarrow \exists x. \ x \in u \land \psi(x, y))$$



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• ψ is function with dom $u \to \psi(u)$ is a set



A10: Axiom of Choice

"Ist T eine Menge, deren sämtliche Elemente yon 0 verschiedene Mengen und untereinander elementenfremd sind, so enthält ihre Yereinigung $\bigcup T$ mindestens eine Untermenge S_1 , welche mit jedem Elemente yon T ein und nur ein Element gemein hat." [Zermelo, 1907]

$$\forall x (x \in z \to x \neq \emptyset \land \\ \forall y (y \in z \to x \cap y = \emptyset \lor x = y))$$

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"additional" axiom



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- "additional" axiom
- ZF versus ZFC

Conservative Extensions



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$$\exists y. \forall x. \neg (x \in y)$$

new symbol $\emptyset \stackrel{\mathsf{cons}}{\Longrightarrow} \forall x. \neg x \in \emptyset$

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$$\forall z_1, z_2. \exists y. \forall x. (x \in y \leftrightarrow x = z_1 \lor x = z_2)$$

new symbol $\{\cdot, \cdot\} \stackrel{\mathsf{cons}\ ex}{\Longrightarrow} \forall z_1, z_2, x. (x \in \{z_1, z_2\} \leftrightarrow x = z_1 \lor x = z_2)$

Conservative Extensions



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A6: Powerset new symbol $\mathbb{P}(\cdot) \stackrel{\mathsf{cons}\ \mathsf{ex}}{\Longrightarrow} \forall x, z. (z \in \mathbb{P}(x) \leftrightarrow \forall u. (u \in z \to u \in x))$



We will use for any formula $\phi(x)$ the syntactical construct

$$\{x \mid \phi(x)\},\$$

called a class term.

Intuitively $\{x \mid \phi(x)\}$ is the collection of all sets a satisfying the formula $\phi(a)$.



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Elimination of class terms:

$$\begin{array}{ll} y \in \{x \mid \phi(x)\} & \text{is replaced by} & \phi(y) \\ \{x \mid \phi(x)\} \in y & \text{is replaced by} & \exists u(u \in y \land \\ & \forall z(z \in u \leftrightarrow \phi(z))) \end{array}$$



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A class term $\{x \mid \phi(x)\}$ does **not** necessarily denote a set.



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Counterexample



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Assume $\{x \mid x \notin x\}$ is a set c, then we obtain

$$c \in c \Leftrightarrow c \notin c$$



A class term $\{x \in A \mid \phi(x)\}$ does denote a set.



A class term $\{x \in A \mid \phi(x)\}$ does denote a set.

Reason



A class term $\{x \in A \mid \phi(x)\}$ does denote a set.

Reason

A3: $\forall x. \exists y. \forall z. (z \in y \leftrightarrow z \in x \land \phi(z))$ for all formulas ϕ



A class term $\{x \in A \mid \phi(x)\}$ does denote a set.

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Class Terms as Subsets



A class term $\{x \in A \mid \phi(x)\}$ does denote a set.

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Conservative extension: $\forall x, z. (z \in F_{\phi}(x) \leftrightarrow z \in x \land \phi(z))$

Different notation: $\forall x, z. (z \in \{t \in x \mid \phi(x)\} \leftrightarrow z \in x \land \phi(z))$



$$\emptyset = \{x \mid x \neq x\}$$



$$\emptyset = \{x \mid x \neq x\}$$

$$\{a, b\} = \{x \mid x = a \lor x = b\}$$



$$\emptyset = \{x \mid x \neq x\}
\{a, b\} = \{x \mid x = a \lor x = b\}
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$$\emptyset = \{x \mid x \neq x\}
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\langle a, b \rangle = \{\{a\}, \{a, b\}\}$$



$$\emptyset = \{x \mid x \neq x\}
\{a, b\} = \{x \mid x = a \lor x = b\}
\{a\} = \{a, a\}
\langle a, b\rangle = \{\{a\}, \{a, b\}\}$$

 $\langle a, b \rangle$ is called the ordered pair of a and b.



The following formulas follow from the ZF axioms

$$\exists x(x=\emptyset)$$

A4 empty set axiom



The following formulas follow from the ZF axioms

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A5 pair axiom



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Special case of pair axiom



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Special case of pair axiom

Special case of pair axiom





$$\exists y \forall z (z \in y \leftrightarrow z \in a \land z \in b)$$

$$y = a \cap b$$



$$\exists y \forall z (z \in y \leftrightarrow z \in a \land z \in b)$$

$$\exists y \forall z (z \in y \leftrightarrow z \in a \lor z \in b)$$

$$v = a \cap b$$

$$y = a \cup b$$



The following theorems are derivable in ZF:

$$\exists y \forall z (z \in y \leftrightarrow z \in a \land z \in b)$$

$$y = a \cap b$$

$$\exists y \forall z (z \in y \leftrightarrow z \in a \lor z \in b)$$

$$y = a \cup b$$

If A is a non-empty class term, then there is a set c satisfying $\forall z (z \in c \leftrightarrow \forall u (u \in A \rightarrow z \in u))$ $c = \bigcap A$



$$\exists y \forall z (z \in y \leftrightarrow z \in a \land z \in b)$$

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$$y = \bigcup a$$



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$$\exists y \forall z (z \in y \leftrightarrow \exists u (u \in a \land z \in u)$$

$$y = \bigcup a$$

$$\blacksquare \exists y \forall z (z \in y \leftrightarrow z \in a \land z \not\in b)$$

$$y = a \setminus b$$



Goal $\exists y \forall z (z \in y \leftrightarrow z \in a \land z \in b)$



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Start with the subset axiom A3

$$\exists y \forall z (z \in y \leftrightarrow z \in x \land \phi(z)).$$



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$$egin{array}{lll} x & \mbox{by} & a \ \phi(z) & \mbox{by} & z \in b \ \mbox{yields} \end{array}$$

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as required

$$a \cap b = \{z \in a \mid z \in b\}$$



Let a, b be sets.

We seek c with $\forall z(z \in c \leftrightarrow (z \in a \lor z \in b))$



Let a, b be sets. We seek c with $\forall z (z \in c \leftrightarrow (z \in a \lor z \in b))$ The pair axioms, A5, $\exists y \forall x (x \in y \leftrightarrow x = z_1 \lor x = z_2)$ guarantees the existence of the set $d = \{a, b\}$



Let a,b be sets. We seek c with $\forall z(z \in c \leftrightarrow (z \in a \lor z \in b))$ The pair axioms, A5, $\exists y \forall x(x \in y \leftrightarrow x = z_1 \lor x = z_2)$ guarantees the existence of the set $d = \{a,b\}$ The sum axiom, A7, $\exists y \forall z(z \in y \leftrightarrow \exists u(z \in u \land u \in x))$ yields the existence of a set c satisfying

$$\forall z(z \in c \leftrightarrow \exists u(z \in u \land u \in d))$$



Let a, b be sets. We seek c with $\forall z(z \in c \leftrightarrow (z \in a \lor z \in b))$ The pair axioms, A5, $\exists y \forall x(x \in y \leftrightarrow x = z_1 \lor x = z_2)$ guarantees the existence of the set $d = \{a, b\}$ The sum axiom, A7, $\exists y \forall z(z \in y \leftrightarrow \exists u(z \in u \land u \in x))$ yields the existence of a set c satisfying

$$\forall z(z \in c \leftrightarrow \exists u(z \in u \land u \in d))$$

Substituting $d = \{a, b\}$ yields the claim.



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$



The following formula can be proved in ZF:

$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

$\begin{array}{l} \mathsf{Proof} \\ \langle \mathsf{a}_1, \mathsf{a}_2 \rangle = \langle \mathsf{b}_1, \mathsf{b}_2 \rangle \end{array}$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

$$\langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcap \langle a_1, a_2 \rangle = \bigcap \bigcap \langle b_1, b_2 \rangle$$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

Proof
$$\langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcap \langle a_1, a_2 \rangle = \bigcap \langle b_1, b_2 \rangle$$

$$\Rightarrow \quad \bigcap \{\{a_1\}, \{a_1, a_2\}\} = \bigcap \{\{b_1\}, \{b_1, b_2\}\}$$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

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$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

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$$\langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcap \langle a_1, a_2 \rangle = \bigcap \langle b_1, b_2 \rangle$$

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$$\Rightarrow \quad \bigcap (\{a_1\} \cap \{a_1, a_2\}) = \bigcap (\{b_1\} \cap \{b_1, b_2\})$$

$$\Rightarrow \quad \bigcap \{a_1\}) = \bigcap \{b_1\}$$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

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 \Rightarrow $\bigcap \{\{a_1\}, \{a_1, a_2\}\} = \bigcap \{\{b_1\}, \{b_1, b_2\}\}$
 \Rightarrow $\bigcap (\{a_1\} \cap \{a_1, a_2\}) = \bigcap (\{b_1\} \cap \{b_1, b_2\})$
 \Rightarrow $\bigcap \{a_1\} = \bigcap \{b_1\}$
 \Rightarrow $a_1 = b_1$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

Case
$$a_1 = a_2$$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

$$\begin{array}{ll} \mathsf{Proof} \\ \langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle & \Rightarrow & \bigcap \langle a_1, a_2 \rangle = \bigcap \langle b_1, b_2 \rangle \\ & \Rightarrow & \bigcap \{ \{a_1\}, \{a_1, a_2\} \} = \bigcap \{ \{b_1\}, \{b_1, b_2\} \} \\ & \Rightarrow & \bigcap (\{a_1\} \cap \{a_1, a_2\}) = \bigcap (\{b_1\} \cap \{b_1, b_2\}) \\ & \Rightarrow & \bigcap \{a_1\}) = \bigcap \{b_1\} \\ & \Rightarrow & a_1 = b_1 \\ \mathsf{Case} \ a_1 = a_2 & \mathsf{Note} \ a_1 = a_2 \Leftrightarrow b_1 = b_2 \\ \end{array}$$



$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

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$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

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Ordered Pairs



The following formula can be proved in ZF:

$$\forall x_1, x_2, y_1, y_2 (\langle x_1, x_2 \rangle = \langle y_1, y_2 \rangle \leftrightarrow x_1 = y_1 \land x_2 = y_2)$$

$$Proof \langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcap \langle a_1, a_2 \rangle = \bigcap \langle b_1, b_2 \rangle \\ \quad \Rightarrow \quad \bigcap \{\{a_1\}, \{a_1, a_2\}\} = \bigcap \{\{b_1\}, \{b_1, b_2\}\} \} \\ \quad \Rightarrow \quad \bigcap \{\{a_1\} \cap \{a_1, a_2\}\} = \bigcap (\{b_1\} \cap \{b_1, b_2\}) \} \\ \quad \Rightarrow \quad \bigcap \{a_1\} \cap \{a_1, a_2\} \cap \{b_1\} \} \\ \quad \Rightarrow \quad a_1 = b_1$$

$$Case \ a_1 = a_2 \qquad \qquad Note \ a_1 = a_2 \Leftrightarrow b_1 = b_2$$

$$Case \ a_1 \neq a_2 \qquad \langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcup (\bigcup \langle a_1, a_2 \rangle \setminus \bigcap \langle a_1, a_2 \rangle) = \bigcup (\bigcup \langle b_1, b_2 \rangle \setminus \bigcap \langle b_1, b_2 \rangle \cap \{b_1, b_2 \rangle \cap \{b_1, b_2 \rangle \cap \{b_1, b_2 \} \cap \{b_2, b_2 \} \cap \{b_1, b_2 \} \cap \{b_2, b_2 \} \cap \{b_1, b_2 \} \cap \{b_2, b_2 \} \cap \{b_2, b_2 \} \cap \{b_1, b_2 \} \cap \{b_2, b_2 \} \cap$$

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$$Case \ a_1 \neq a_2 \\ \langle a_1, a_2 \rangle = \langle b_1, b_2 \rangle \quad \Rightarrow \quad \bigcup (\bigcup \langle a_1, a_2 \rangle \setminus \bigcap \langle a_1, a_2 \rangle) = \bigcup (\bigcup \langle b_1, b_2 \rangle \setminus \bigcap \langle b_1, b_2 \rangle \\ \quad \Rightarrow \quad \bigcup \{a_1, a_2\} \setminus \{a_1\} = \bigcup \{b_1, b_2\} \setminus \{b_1\} \}$$

$$\Rightarrow \quad \bigcup \{a_2\} = \bigcup \{b_2\}$$

Ordered Pairs



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$$Case \ a_1 = a_2 \qquad \qquad Note \ a_1 = a_2 \Leftrightarrow b_1 = b_2$$

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• A relation r is a set of ordered pairs, i.e.

$$rel(r) \equiv \forall x (x \in r \rightarrow \exists x_1, x_2 (x = \langle x_1, x_2 \rangle))$$



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■ The relation r is said to be a relation on the set s if

$$rel(r,s) \equiv rel(r) \land \forall x_1, x_2 (\langle x_1, x_2 \rangle \in r \rightarrow x_1 \in s \land x_2 \in s)$$



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A (weak) function is a one-valued relation, i.e.

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- A (weak) function is a one-valued relation, i.e.

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A function f is said to be a function from a set a to a set b if $func(f, a, b) \equiv func(f) \land \forall x_1, x_2(\langle x_1, x_2 \rangle \in f \rightarrow x_1 \in a \land x_2 \in b)$





From the ZF axioms we can prove for any sets a, b the existence

• of the set of all relations on a



- of the set of all relations on a
- of the set of all functions from a to b



- of the set of all relations on a
- of the set of all functions from a to b
- i.e.



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- i.e.
- $\forall x \exists y \forall z (z \in y \leftrightarrow rel(z, x))$



- of the set of all relations on a
- of the set of all functions from a to b
- i.e.
- $\forall x \exists y \forall z (z \in y \leftrightarrow rel(z, x))$
- $\forall u, w \exists y \forall z (z \in y \leftrightarrow func(z, u, w))$



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- not universally accepted
 (e.g., provides a handle on objects of which only existence is known)



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 - ...

Towards Macro Structures

Natural Numbers $\mathbb N$



Define for any set a its successor set a^+ :

$$a^+ = a \cup \{a\}$$



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$$\{\emptyset,\emptyset^+,\emptyset^{++},\ldots\}$$



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$$2=1^+=\{\emptyset,\{\emptyset\}\}$$



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$$\{\emptyset, \emptyset^+, \emptyset^{++}, \ldots\}$$

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$$1 = \emptyset^+ = \{\emptyset\}$$

$$2 = 1^+ = \{\emptyset, \{\emptyset\}\}$$

$$3 = 2^+ = \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}\}\}$$

Formal Definition of $\mathbb N$



$$Ded(a) \equiv 0 \in a \land \forall x (x \in a \rightarrow x^+ \in a)$$

Formal Definition of $\mathbb N$



$$Ded(a) \equiv 0 \in a \land \forall x (x \in a \rightarrow x^+ \in a)$$

a is called a Dedekind set or inductive if Ded(a) is true.

Formal Definition of $\mathbb N$



$$Ded(a) \equiv 0 \in a \land \forall x (x \in a \rightarrow x^+ \in a)$$

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$$\mathbb{N} = \bigcap \{a \mid Ded(a)\}$$



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Peano's Axioms



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 $\forall n, m (n \in \mathbb{N} \land m \in \mathbb{N} \land n^+ = m^+ \rightarrow n = m).$



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The foundation axiom, A2,

$$\exists y(y \in x) \to \exists y(y \in x \land \forall z \neg (z \in x \land z \in y)),$$

instantiated for $x = \{n, m\}$ yields

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Thus $n \notin m$ or $m \notin n$.



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Case n=0

By definition $n^+ = \{0\}$.

Thus obviously $0 \in n^+$ and also $n^+ \in x$.

Transitive Sets



Definition

A set a is called transitive if every element of a is also a subset of a. In symbols

$$trans(a) \leftrightarrow \forall x (x \in a \rightarrow x \subseteq a)$$

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Lemma

- ① n is transitive for all $n \in \mathbb{N}$.



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If $x \in n$ then by hypothesis $x \subseteq n \subseteq n^+$.

If x = n, then by definition $x \subseteq n^+$.

$\mathbb N$ is transitive



Prove $\forall n (n \in \mathbb{N} \to n \subseteq \mathbb{N})$ by induction.

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For n = 0 this is clear.

If $n \in \mathbb{N}$ and by induction hypothesis $n \subseteq \mathbb{N}$

then also $n^+ = n \cup \{n\} \subseteq \mathbb{N}$.



Claim

The \in -relation is the smallest transitive relation r on $\mathbb N$ with $\langle n, n^+ \rangle \in r$ for all n. i.e.

$$\forall n, m(n \in m \to \langle n, m \rangle \in r)$$



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Case n = m We immediately have $\langle m, m^+ \rangle \in r$.



The <-relation on $\mathbb N$ coincides with the \in -relation.



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Any natural number n is the set of all its predecessors, i.e.

$$n = \{m \mid m < n\}.$$

The Recursion Theorem



Let F be a function satisfying $rng(F) \subseteq dom(F)$ and let u be an element in dom(F).

Then there exists exactly one function f satisfying

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- **2** f(0) = u,

The assumptions $rng(F) \subseteq dom(F)$ and $u \in dom(F)$ are needed to make sure that all function applications of F are defined.



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Since
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$$n \in x$$
 implies $n^+ \in x$
since $f(n^+) = F(f(n)) = F(g(n)) = g(n^+)$



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Thus by the last Peano axiom induction axiom, we get

$$x = \mathbb{N}$$

i.e.
$$f = g$$
.

Idea of Existence Proof



Idea

$$H = \{h \mid func(h) \wedge h(0) = u \wedge \exists n(n \neq 0 \wedge dom(h) = n \\ \wedge \forall m(m^+ \in n \rightarrow h(m^+) = F(h(m)))\}$$

Idea of Existence Proof



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$$H = \{h \mid func(h) \land h(0) = u \land \exists n(n \neq 0 \land dom(h) = n \land \forall m(m^+ \in n \rightarrow h(m^+) = F(h(m))))\}$$

and

$$f = \bigcup H$$

Addition



for every $m \in \mathbb{N}$ there is a unique function add_m such that

$$\begin{array}{lcl} add_m(0) & = & m \\ add_m(n^+) & = & (add_m(n))^+ \end{array}$$

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for every $m \in \mathbb{N}$ there is a unique function add_m such that

$$add_m(0) = m$$

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Apply the recursion theorem with u = m and $F(x) = x^+$

Multiplication



for every $m \in \mathbb{N}$ there is a unique function $mult_m$ such that

$$mult_m(0) = 0$$

 $mult_m(n^+) = add_m(mult_m(n))$

Multiplication



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Apply the recursion theorem with u = 0 and $F_m(x) = add_m(x)$.

The Integers



The idea is to reconstruct an integer

Z

as a pair $\langle m, n \rangle$ of natural numbers with

$$z = m - n$$

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The Integers



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as a pair $\langle m, n \rangle$ of natural numbers with

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Since $\langle 5,7\rangle$ and $\langle 8,10\rangle$ would both represent the same number, we have to use equivalence classes of ordered pairs instead of pairs themselves.



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An ordinal x is a set such that (x, \in) is a well-ordered set.



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Definition

A set a is called an ordinal if

- ① it is transitive, and
- it is totally ordered by inclusion

(equivalently: a transitive set of transitive sets)

We will denote ordinals by lowercase Greek letter, α , β ,



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- 3 The set of all natural numbers, traditionally denoted by the letter ω , is an ordinal.
- **4** If α is an ordinal, then every element $\beta \in \alpha$ is an ordinal.

Two Types of Ordinals



① An ordinal α such that $\alpha = \beta^+ = \beta \cup \{\beta\}$ for some β is called a *successor ordinal*.

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① An ordinal α such that $\alpha = \beta^+ = \beta \cup \{\beta\}$ for some β is called a *successor ordinal*.

② An ordinal α such that for all β with $\beta \in \alpha$ there is γ such that $\beta \in \gamma$ and $\gamma \in \alpha$ is called a *limit ordinal*.

Representation Theorem



For every well-ordered set (G, <) there is a unique ordinal α such that

$$(G,<)\cong(\alpha,\epsilon)$$



1 0,
$$0^+ = 1$$
, $0^{++} = 2$, ..., n , ...



① 0,
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, $0^{++} = 2$, ..., n , ...

$$\bullet$$
 ω , $\omega + 1$, $\omega + 2$, $\ldots \omega + n \ldots$



- **1** 0, $0^+ = 1$, $0^{++} = 2$, ..., n, ...
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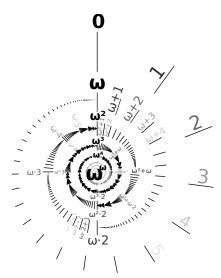
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Visualisation





[found on wikipedia]

Limitations of ZFC



Gödel's Second Incompleteness Theorem

Assume T is a consistent theory which contains elementary arithmetic. Then $T \not\vdash Cons(T)$; the consistency of T cannot be proved from T.

Continuuom Hypothesis independent from ZFC

There is no set whose cardinality lies strictly between that of the integers and the reals.