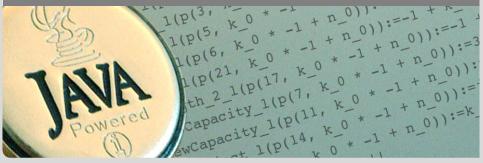


# Applications of Formal Verification Functional Verification of Java Programs: Java Modelling Language

Prof. Dr. Bernhard Beckert · Dr. Vladimir Klebanov | SS 2012

KIT – INSTITUT FÜR THEORETISCHE INFORMATIK





### Idea

Specifications fix a contract between caller and callee of a method (between client and implementor of a module):

If caller guarantees precondition then callee guarantees certain outcome

- Interface documentation
- Contracts described in a mathematically precise language (JML)
  - higher degree of precision
  - automation of program analysis of various kinds (runtime assertion checking, static verification)
- Note: Errors in specifications are at least as common as errors in code,



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#### /\*@ public normal\_behavior

- @ requires pin == correctPin;
- @ ensures customerAuthenticated;

```
@*/
```

### public void enterPIN (int pin) {

• • •

- Java comments with '@' as first character are JML specifications
- Within a JML annotation, an '@' is ignored:
  - if it is the first (non-white) character in the line
  - if it is the last character before '\*/'.

 $\Rightarrow$  The blue '@'s are not required, but it's a *convention* to use them.



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```
/*@ public normal_behavior //<hello!<
    @ requires pin == correctPin;
    @ ensures customerAuthenticated;
    @*/
public void enterPIN (int pin) {
    ...</pre>
```

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```
public class ATM {
    private /*@ spec_public @*/ BankCard insertedCard = null;
    private /*@ spec_public @*/
        boolean customerAuthenticated = false;
    /*@ public normal behavior ... @*/
```

- Modifiers to specification cases have no influence on their semantics.
- *public* specification items cannot refer to *private* fields.
- Private fields can be declared public for specification purposes only.



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```
/*@ requires r;
```

- @ assignable a;
- @ diverges d;
- @ ensures post;
- @ signals\_only E1, ..., En;
- @ signals(E e) s;

```
@*/
```

```
T m(...);
```

#### Abbreviations

```
normal behavior = signals(Exception e) false;
```

exceptional\_behavior = ensures false;

#### keyword 'also' separates the contracts of a method

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```
/*@ requires r; //what is the caller's obligation?
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@ <b>assignable</b> a;	<pre>//what is the caller's obligation? //which locations may be assigned by m //when may m non-terminate?,En;</pre>	?
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@ <b>signals_only</b> E1,, En; //what exc-types may be thrown?
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### keyword 'also' separates the contracts of a method



- can be placed anywhere in a class (or interface)
- express global consistency properties (not specific to a particular method)
- must hold "always" (cf. visible state semantics, observed state semantics)
- instance invariants can, static invariants cannot refer to this
- default: instance within classes, static within interfaces



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## **Pure Methods**



### Pure methods terminate and have no side effects.

After declaring

```
public /*@ pure @*/ boolean cardIsInserted() {
    return insertedCard!=null;
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cardIsInserted()

could replace

insertedCard != null

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#### in JML annotations.



#### 'pure' $\approx$ 'diverges false; ' + 'assignable \nothing; '

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### Expressions



#### All Java expressions without side-effects

- ==>, <==>: implication, equivalence
- I \forall, \exists
- $\texttt{ \num_of, \sum, \product, \min, \max} \\$
- \old(...): referring to pre-state in postconditions
- \result: referring to return value in postconditions

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(\forall int i; 0<=i && i<\result.length; \result[i]>0)
equivalent to
(\forall int i; 0<=i && i<\result.length ==> \result[i]>0)
(\exists int i; 0<=i && i<\result.length; \result[i]>0)
equivalent to
(\exists int i; 0<=i && i<\result.length && \result[i]>0)

- Note that quantifiers bind two expressions, the range predicate and the body expression.
- A missing range predicate is by default true.
- JML excludes null from the range of quantification.



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# Generalised and Numerical Quantifiers



(\num_of T i; e)	ł
(\sum T i; p; t)	
(\product T i; p; t)	
(\min T i; p; t)	
(\max T i; p; t)	

 $\#\{i|[e]\},$  number of elements of type T with property e

$$\sum_{i:[\rho]} [t]$$

$$\prod_{i:[\rho]} [t]$$

$$min\{[t]\}$$

$$max_{i:[\rho]}$$

$$i:[\rho]$$



Comma-separated list of:

• e.f (where f a field)



Comma-separated list of:

- e.f (where f a field)
- a[\*], a[x..y] (where a an array expression)



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C x, y; int i;
//@ assignable x, x.i;
void m() {
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Comma-separated list of:

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```
C x, y; int i;
//@ assignable x, x.i;
void m() {
    C tmp = x;
    tmp.i = 27;
    x = y;
    x.i = 27;
}
```



Comma-separated list of:

- e.f (where f a field)
- a[\*], a[x..y] (where a an array expression)
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C x, y; int i;
//@ assignable x, x.i;
void m() {
    C tmp = x; //allowed (local variable)
    tmp.i = 27;
    x = y;
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Comma-separated list of:

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C x, y; int i;
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  tmp.i = 27; //allowed (in assignable clause)
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   x.i = 27; //forbidden (not local, not in assignable)
}
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#### Example

C x, y; int i; //@ assignable x, x.i; void m() { C tmp = x; //allowed (local variable) tmp.i = 27; //allowed (in assignable clause) x = y; //allowed (in assignable clause) x.i = 27; //forbidden (not local, not in assignable) }

#### assignable clauses are always evaluated in the pre-state!



# with a boolean JML expression e specifies that the method may may not terminate only when e is true in the pre-state.

#### Examples

diverges false; The method must always terminate. diverges true; The method may terminate or not.



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diverges false; The method must always terminate. diverges true; The method may terminate or not.



```
ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
signals (En en) sn;
```

- normal termination  $\Rightarrow$  p must hold (in post-state)
- exception thrown  $\Rightarrow$  must be of type ET1, ..., or ETm
- exception of type E1 thrown  $\Rightarrow$  s1 must hold (in post-state)
- exception of type En thrown ⇒ sn must hold (in post-state)



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ensures p;
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#### • normal termination $\Rightarrow$ p must hold (in post-state)

- exception thrown  $\Rightarrow$  must be of type ET1, ..., or ETm
- exception of type E1 thrown  $\Rightarrow$  s1 must hold (in post-state)

# ■ exception of type En thrown ⇒ sn must hold (in post-state)



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- exception of type E1 thrown  $\Rightarrow$  s1 must hold (in post-state)
- exception of type En thrown ⇒ sn must hold (in post-state)



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ensures p;
signals_only ET1, ..., ETm;
signals (E1 e1) s1;
...
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public interface IBonusCard {

#### public void addBonus(int newBonusPoints);



}



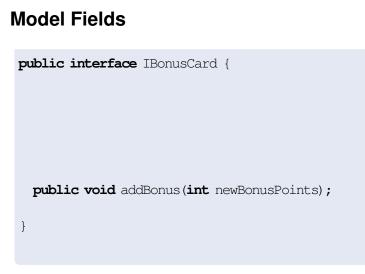
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public interface IBonusCard {

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/*@ public instance model int bonusPoints; @*/
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### **Model Fields**

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/\*@ ensures bonusPoints == \old(bonusPoints) +newBonusPoints;

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### **Model Fields**



public interface IBonusCard {

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```
/*@ ensures bonusPoints == \old(bonusPoints) +newBonusPoints;
@ assignable bonusPoints;
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```

## Implementing Interfaces



### public interface IBonusCard {

/\*@ public instance model int bonusPoints; @\*/

```
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#### Implementation

public class BankCard implements IBonusCard{
 public int bankCardPoints;
/\*@ private represents bonusPoints = bankCardPoints; @\*/
 public void addBonus(int newBonusPoints) {
 bankCardPoints += newBonusPoints; }
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## **Other Representations**



/\*@ represents x \such\_that A(x); @\*/

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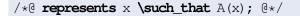


## **Other Representations**



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### Inheritance of Specifications in JML



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 An operation contract is inherited by all overridden methods.

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- loop invariants '//@ loop\_invariant p;'
- data groups
- refines
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#### JML has modifiers non\_null and nullable

private /\*@spec\_public non\_null@\*/ Object x;

wimplicit invariant added to class: 'invariant x != null;'

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void m(/*@non_null@*/ Object p);
```

wimplicit precondition added to all contracts:
'requires p != null;'

/\*@**non\_null**@\*/ Object m();

wimplicit postcondition added to all contracts:
 'ensures \result != null;'

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If something may be null, you have to declare it nullable

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/*@ requires y >= 0;
@ ensures \result >= 0;
@ ensures \result * \result <= y;
@ ensures (\result+1) * (\result+1) > y;
@ */
public static int isqrt(int y)
```

For y = 1 and  $\result = 1073741821 = \frac{1}{2}(MAX_INT - 5)$  the above postcondition is true, though we do not want 1073741821 to be a square root of 1. JML uses the Java semantics of integers:

1073741821 \* 1073741821 = -21474836391073741822 \* 1073741822 = 4

The JML type \bigint provides arbitrary precision integers.

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#### Many tools support JML (see JML homepage). Among them:

- jml: JML syntax checker
- jmldoc: code documentation (like Javadoc)
- jmlc: compiles Java+JML into bytecode with assertion checks
- jmlunit: unit testing (like JUnit)
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