KI-Programmierung

Basic Search Algorithms

Bernhard Beckert

UNIVERSITÄT KOBLENZ-LANDAU

Winter Term 2007/2008

Example: Travelling in Romania

Scenario

On holiday in Romania; currently in Arad

Flight leaves tomorrow from Bucharest

Example: Travelling in Romania

Scenario

On holiday in Romania; currently in Arad

Flight leaves tomorrow from Bucharest

Goal

Be in Bucharest

Example: Travelling in Romania

Scenario

On holiday in Romania; currently in Arad
Flight leaves tomorrow from Bucharest

Goal

Be in Bucharest

Formulate problem

***States*: various cities**
***Actions*: drive between cities**

Example: Travelling in Romania

Scenario

On holiday in Romania; currently in Arad
Flight leaves tomorrow from Bucharest

Goal

Be in Bucharest

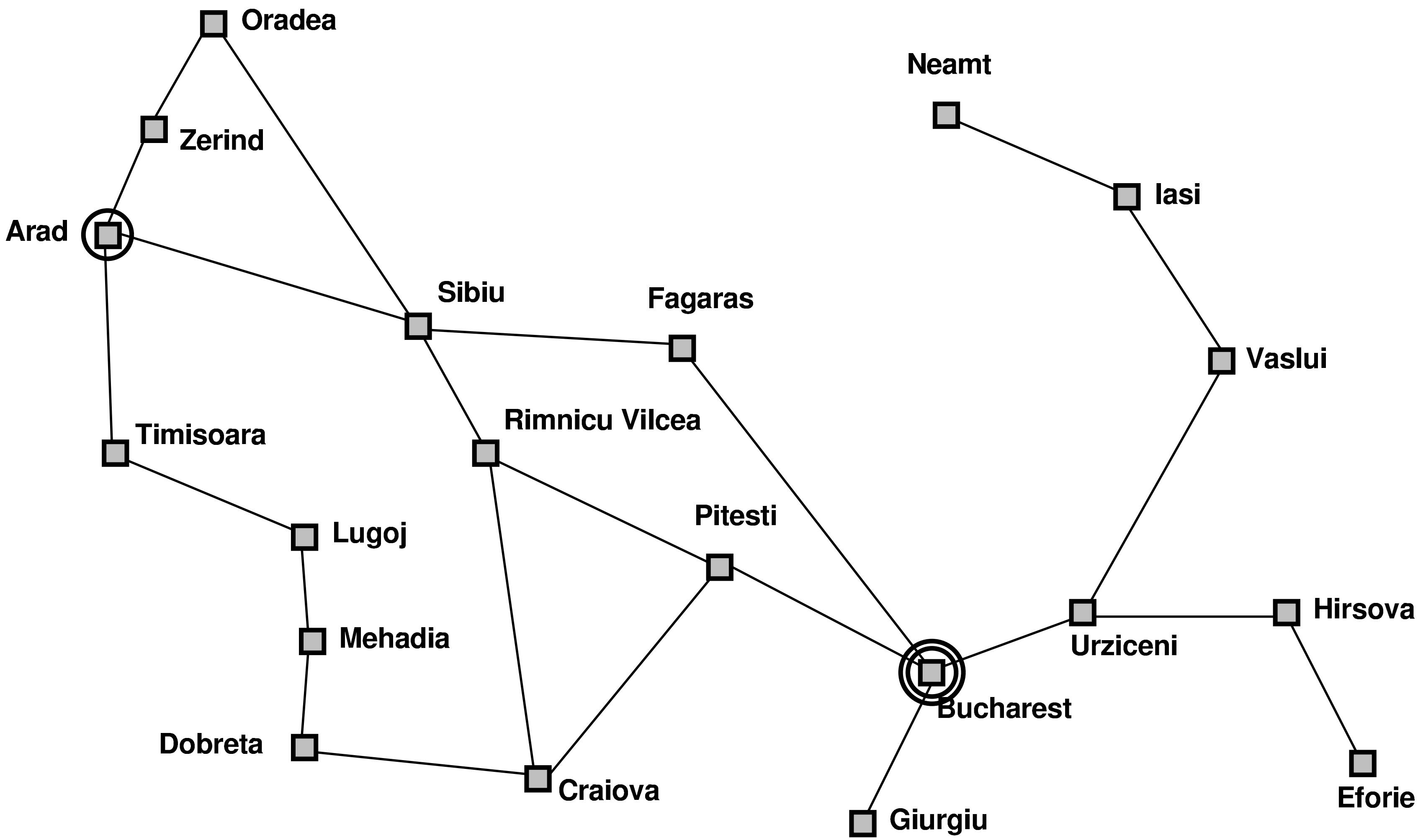
Formulate problem

***States*: various cities**
***Actions*: drive between cities**

Solution

Appropriate sequence of cities
e.g.: Arad, Sibiu, Fagaras, Bucharest

Example: Travelling in Romania



Problem formulation

Defined by the following four items

1. **Initial state**

   Example: ***Arad***

2. **Successor function** S

   Example: $S(\textit{\textbf{Arad}}) = \{ \langle \textit{\textbf{goZerind}}, \textit{\textbf{Zerind}} \rangle, \langle \textit{\textbf{goSibiu}}, \textit{\textbf{Sibiu}} \rangle, \ldots \}$

3. **Goal test**

   Example: $x = \textit{\textbf{Bucharest}}$ **(explicit test)**

4. **Path cost** **(optional)**

   Example: **sum of distances, number of operators executed, etc.**

Problem formulation

Solution

A sequence of operators
leading from the initial state to a goal state

Selecting a state space

Abstraction

Real world is absurdly complex

State space must be abstracted for problem solving

(Abstract) state

Set of real states

(Abstract) operator

Complex combination of real actions

Example: *Arad* $\longrightarrow$ *Zerind* represents complex set of possible routes

(Abstract) solution

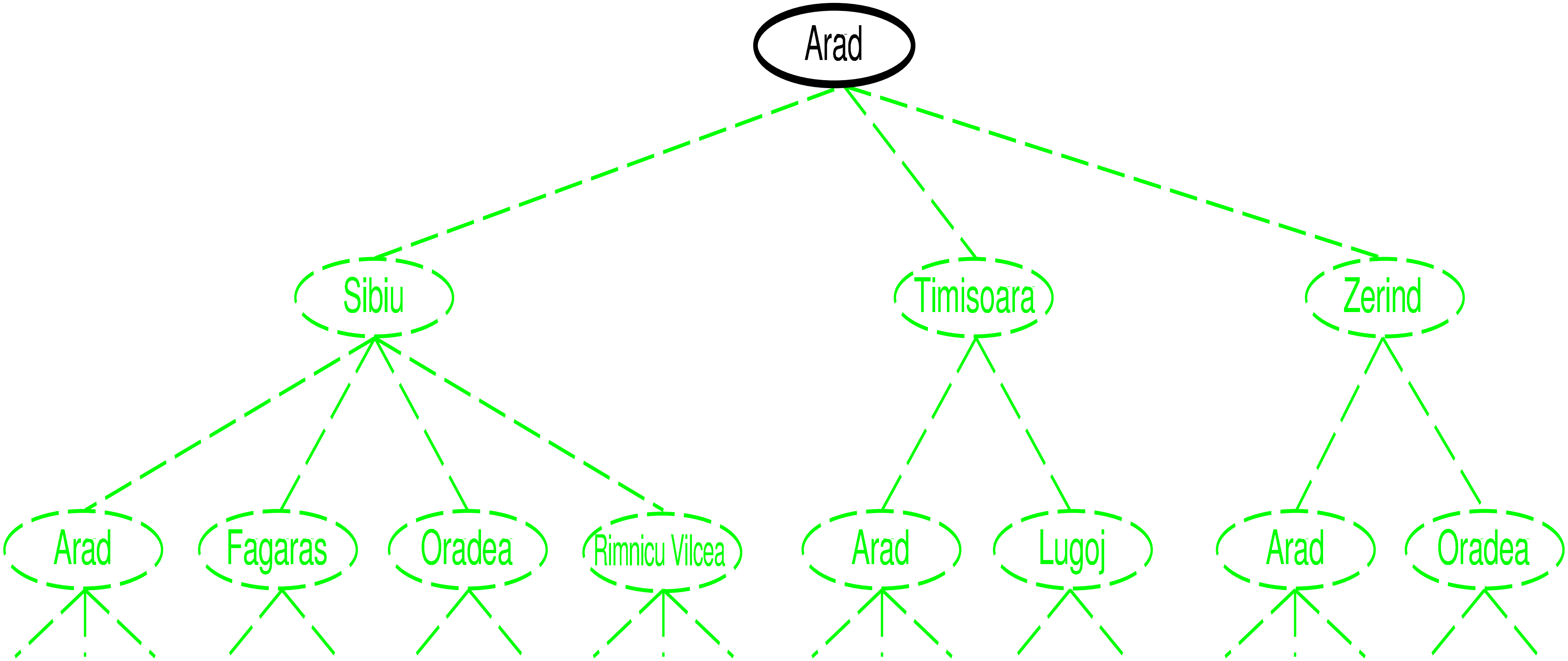
Set of real paths that are solutions in the real world

Tree search algorithms

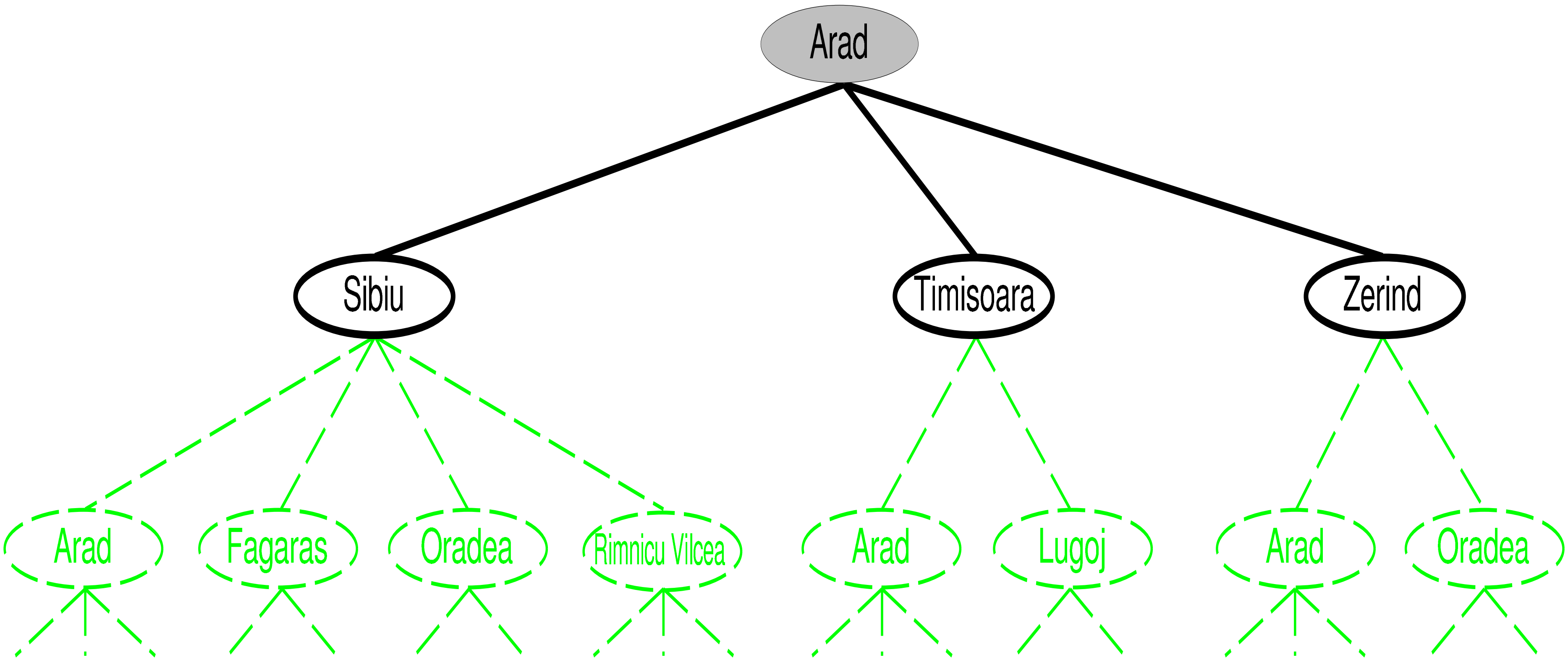
- **Simulated exploration of state space in a search tree by generating successors of already-explored states**

function TREE-SEARCH(*problem*, *strategy*) **returns** a solution or failure
  initialize the search tree using the initial state of *problem*
  loop do
    if there are no candidates for expansion **then return** failure
    choose a leaf node for expansion according to *strategy*
    if the node contains a goal state **then**
      return the corresponding solution
    else
      expand the node and add the resulting nodes to the search tree
  end

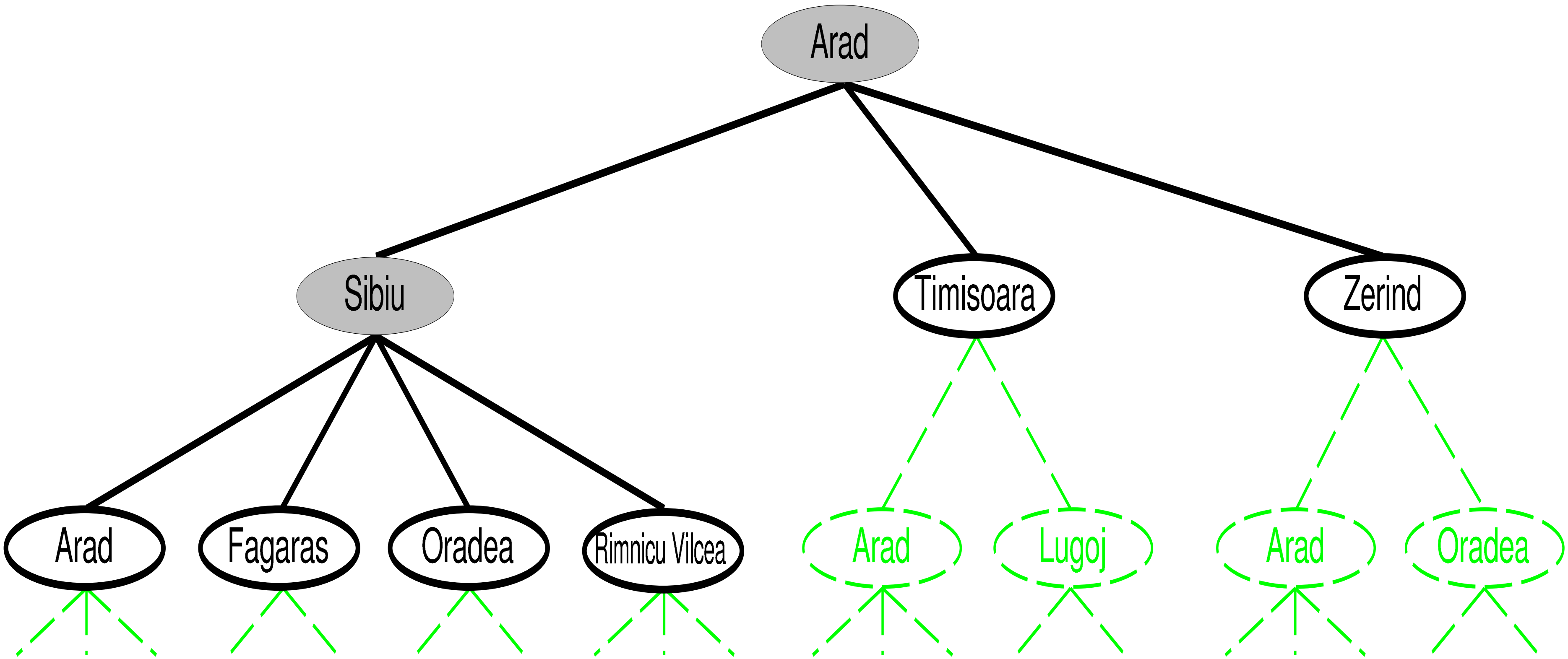
Tree search: Example



Tree search: Example



Tree search: Example



Implementation of search algorithms

function TREE-SEARCH(*problem, fringe*) **returns** a solution or failure

    fringe $\leftarrow$ INSERT(MAKE-NODE(INITIAL-STATE[*problem*]),*fringe*)

    loop do

        if *fringe* is empty **then return** failure

        node $\leftarrow$ REMOVE-FIRST(*fringe*)

        if GOAL-TEST[*problem*] applied to STATE(*node*) succeeds **then**

            return *node*

        else

            fringe $\leftarrow$ INSERT-ALL(EXPAND(*node*, *problem*), *fringe*)

    end

fringe	**queue of nodes not yet considered**
State	**gives the state that is represented by *node***
Expand	**creates new nodes by applying possible actions to *node***

Search strategies

Strategy

Defines the order of node expansion

Important properties of strategies

completeness	**does it always find a solution if one exists?**
time complexity	**number of nodes generated/expanded**
space complexity	**maximum number of nodes in memory**
optimality	**does it always find a least-cost solution?**

Time and space complexity measured in terms of

b	**maximum branching factor of the search tree**
d	**depth of a solution with minimal distance to root**
m	**maximum depth of the state space (may be ∞)**

Uninformed search strategies

Uninformed search

Use only the information available in the problem definition

Frequently used strategies

- **Breadth-first search**
- **Uniform-cost search**
- **Depth-first search**
- **Depth-limited search**
- **Iterative deepening search**

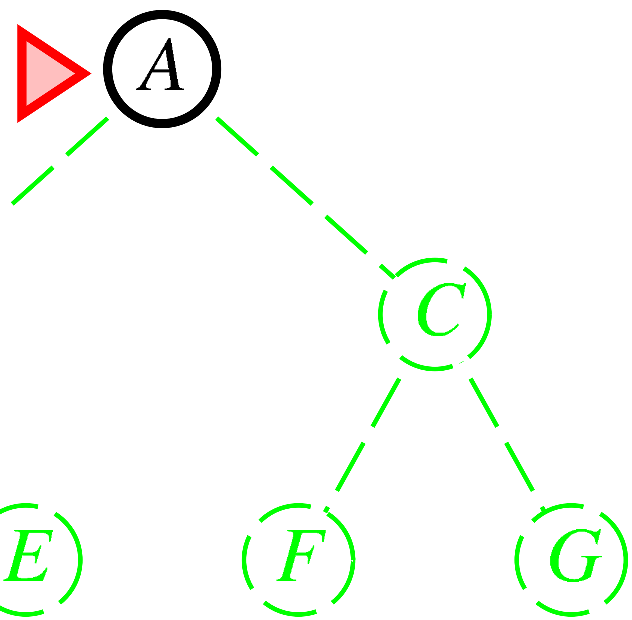
Breadth-first search

Idea

Expand shallowest unexpanded node

Implementation

***fringe* is a FIFO queue, i.e. successors go in at the end of the queue**



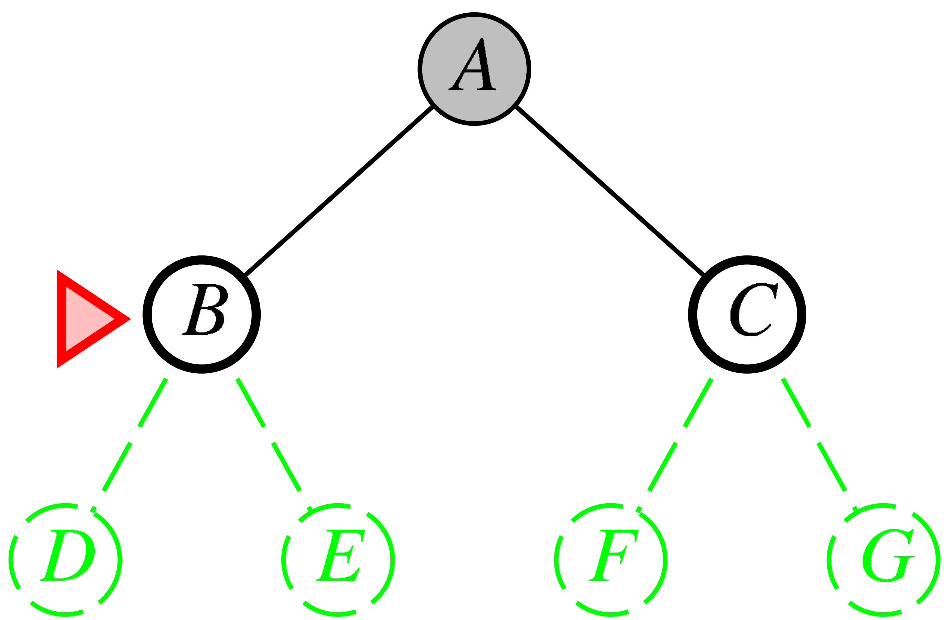
Breadth-first search

Idea

Expand shallowest unexpanded node

Implementation

***fringe* is a FIFO queue, i.e. successors go in at the end of the queue**



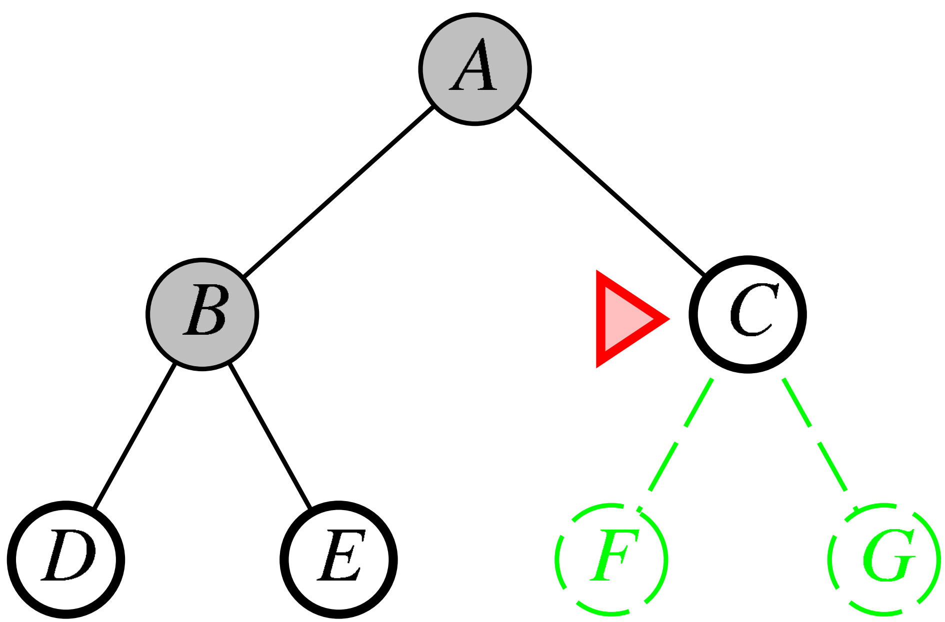
Breadth-first search

Idea

Expand shallowest unexpanded node

Implementation

***fringe* is a FIFO queue, i.e. successors go in at the end of the queue**



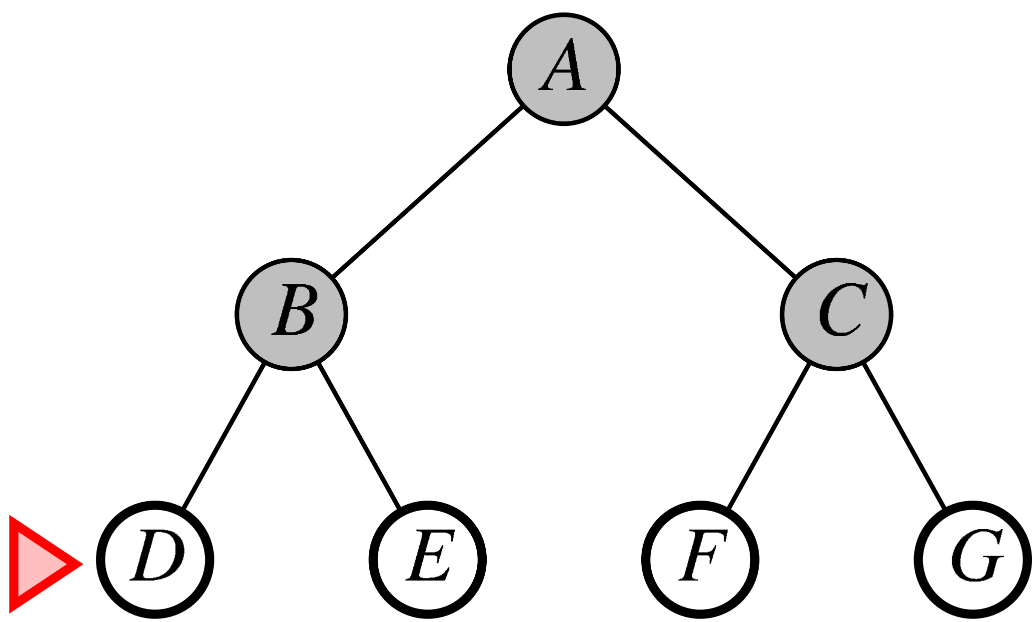
Breadth-first search

Idea

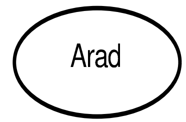
Expand shallowest unexpanded node

Implementation

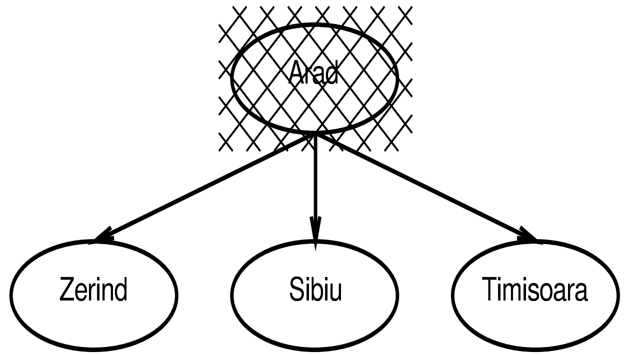
***fringe* is a FIFO queue, i.e. successors go in at the end of the queue**



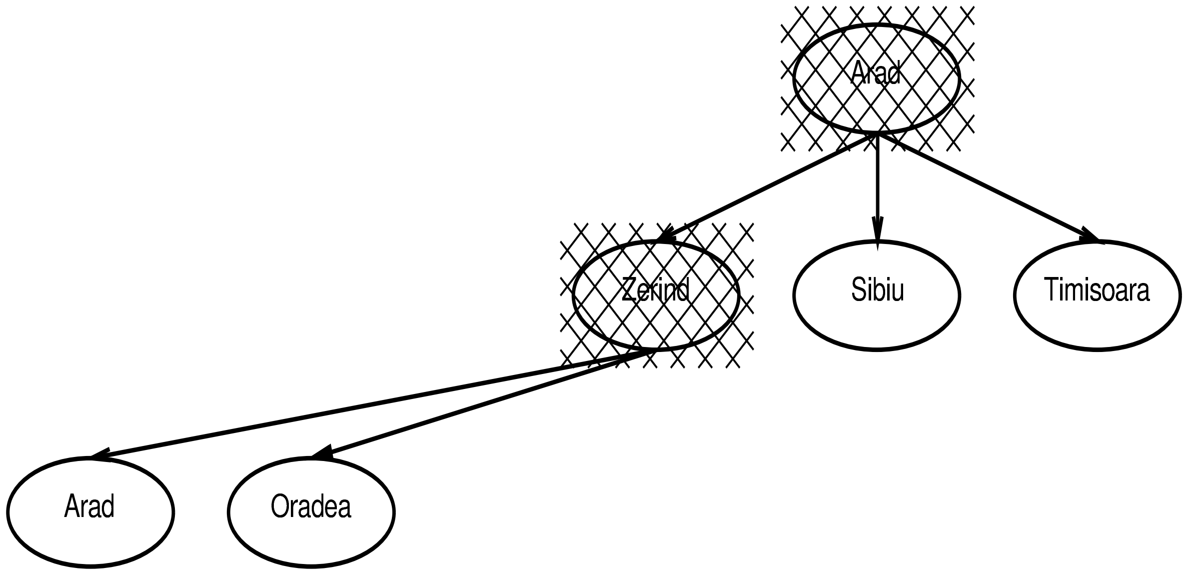
Breadth-first search: Example Romania



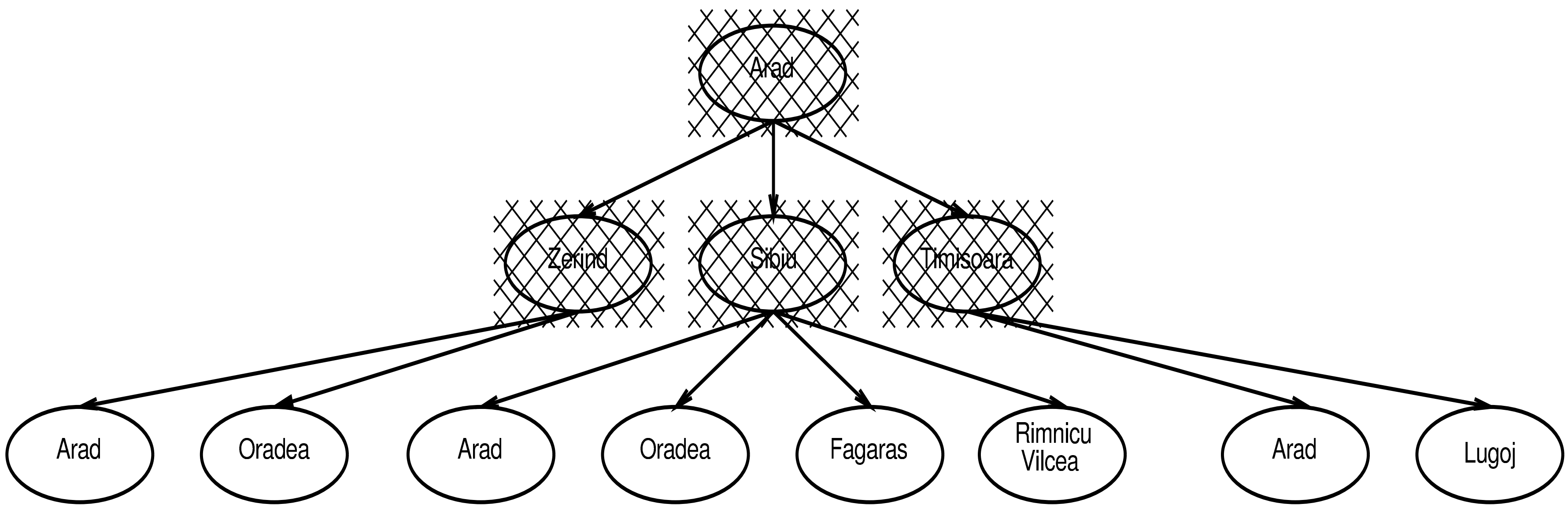
Breadth-first search: Example Romania



Breadth-first search: Example Romania



Breadth-first search: Example Romania



Breadth-first search: Properties

Complete

Time

Space

Optimal

Breadth-first search: Properties

Complete **Yes (if b is finite)**

Time

Space

Optimal

Breadth-first search: Properties

Complete **Yes (if b is finite)**

Time $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) \in O(b^{d+1})$

i.e. exponential in d

Space

Optimal

Breadth-first search: Properties

Complete **Yes (if b is finite)**

Time $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) \in O(b^{d+1})$

i.e. exponential in d

Space $O(b^{d+1})$

keeps every node in memory

Optimal

Breadth-first search: Properties

Complete Yes (if b is finite)

Time $1+b+b^2+b^3+\ldots+b^d+b(b^d-1) \in O(b^{d+1})$

i.e. exponential in d

Space $O(b^{d+1})$

keeps every node in memory

Optimal Yes (if cost = 1 per step), not optimal in general

Breadth-first search: Properties

Complete Yes (if b is finite)

Time $1+b+b^2+b^3+\ldots+b^d+b(b^d-1) \in O(b^{d+1})$

i.e. exponential in d

Space $O(b^{d+1})$

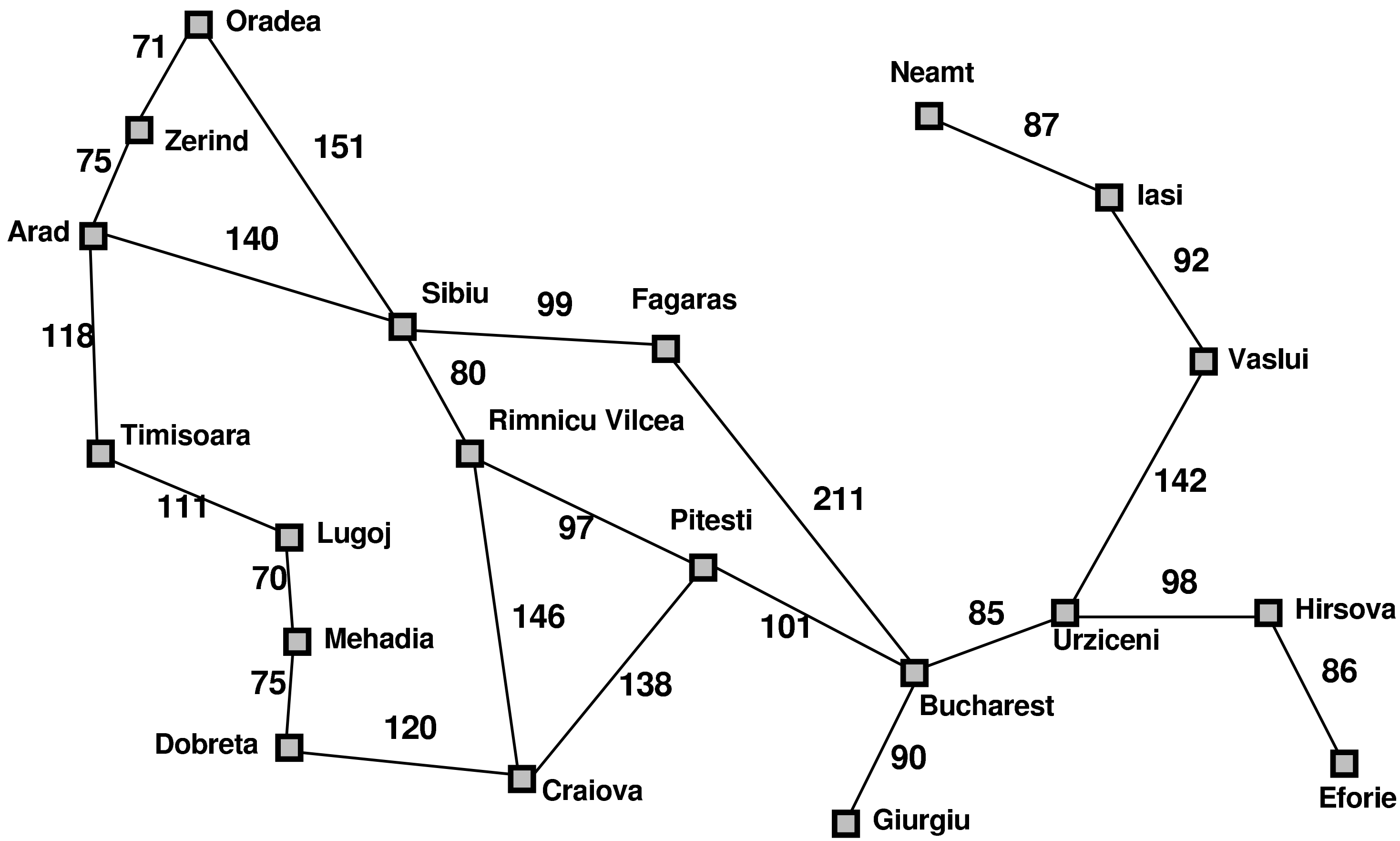
keeps every node in memory

Optimal Yes (if cost = 1 per step), not optimal in general

Disadvantage

Space is the big problem
(can easily generate nodes at 5MB/sec so 24hrs = 430GB)

Romania with step costs in km



Straight–line distance
to Bucharest

Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	178
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	98
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

Uniform-cost search

Idea

Expand least-cost unexpanded node
(costs added up over paths from root to leafs)

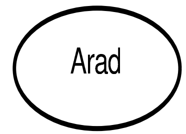
Implementation

***fringe* is queue ordered by increasing path cost**

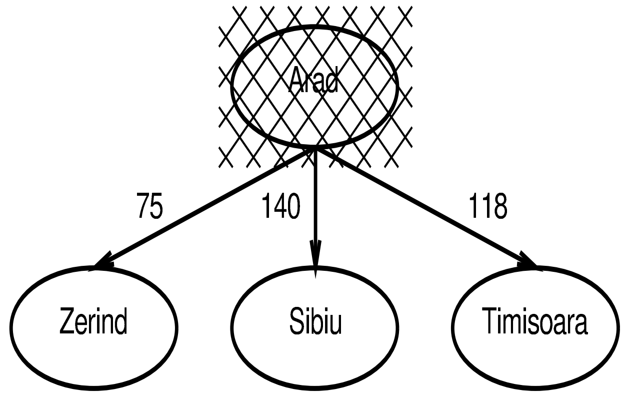
Note

Equivalent to breadth-first search if all step costs are equal

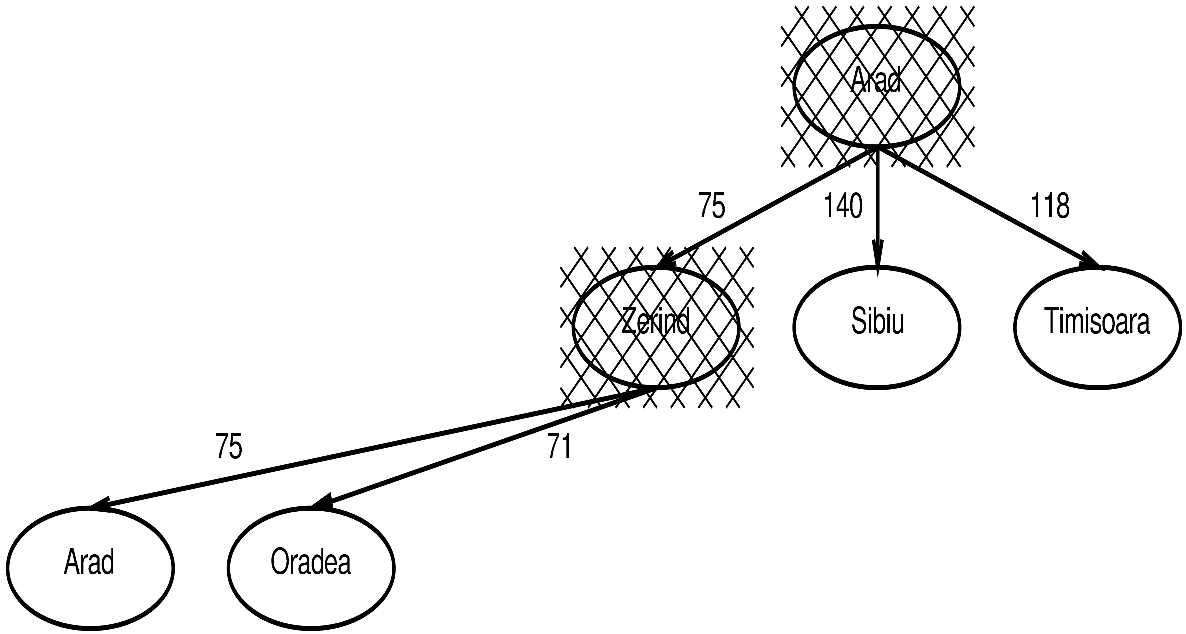
Uniform-cost search



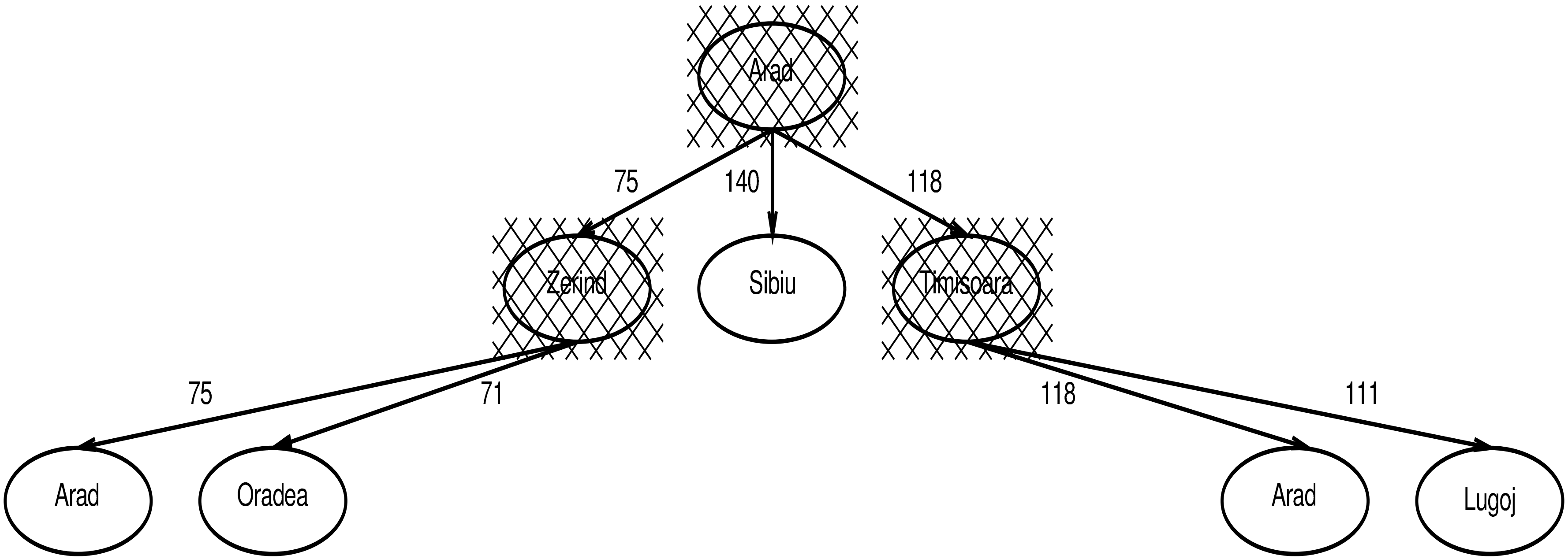
Uniform-cost search



Uniform-cost search



Uniform-cost search



Uniform-cost search: Properties

Complete

Time

Space

Optimal

Uniform-cost search: Properties

Complete Yes (if step costs positive)

Time

Space

Optimal

Uniform-cost search: Properties

Complete	**Yes (if step costs positive)**
Time	**# of nodes with past-cost less than that of optimal solution**
Space	
Optimal	

Uniform-cost search: Properties

Complete	**Yes (if step costs positive)**
Time	**# of nodes with past-cost less than that of optimal solution**
Space	**# of nodes with past-cost less than that of optimal solution**
Optimal	

Uniform-cost search: Properties

Complete	**Yes (if step costs positive)**
Time	**# of nodes with past-cost less than that of optimal solution**
Space	**# of nodes with past-cost less than that of optimal solution**
Optimal	**Yes**

Depth-first search

Idea

Expand deepest unexpanded node

Implementation

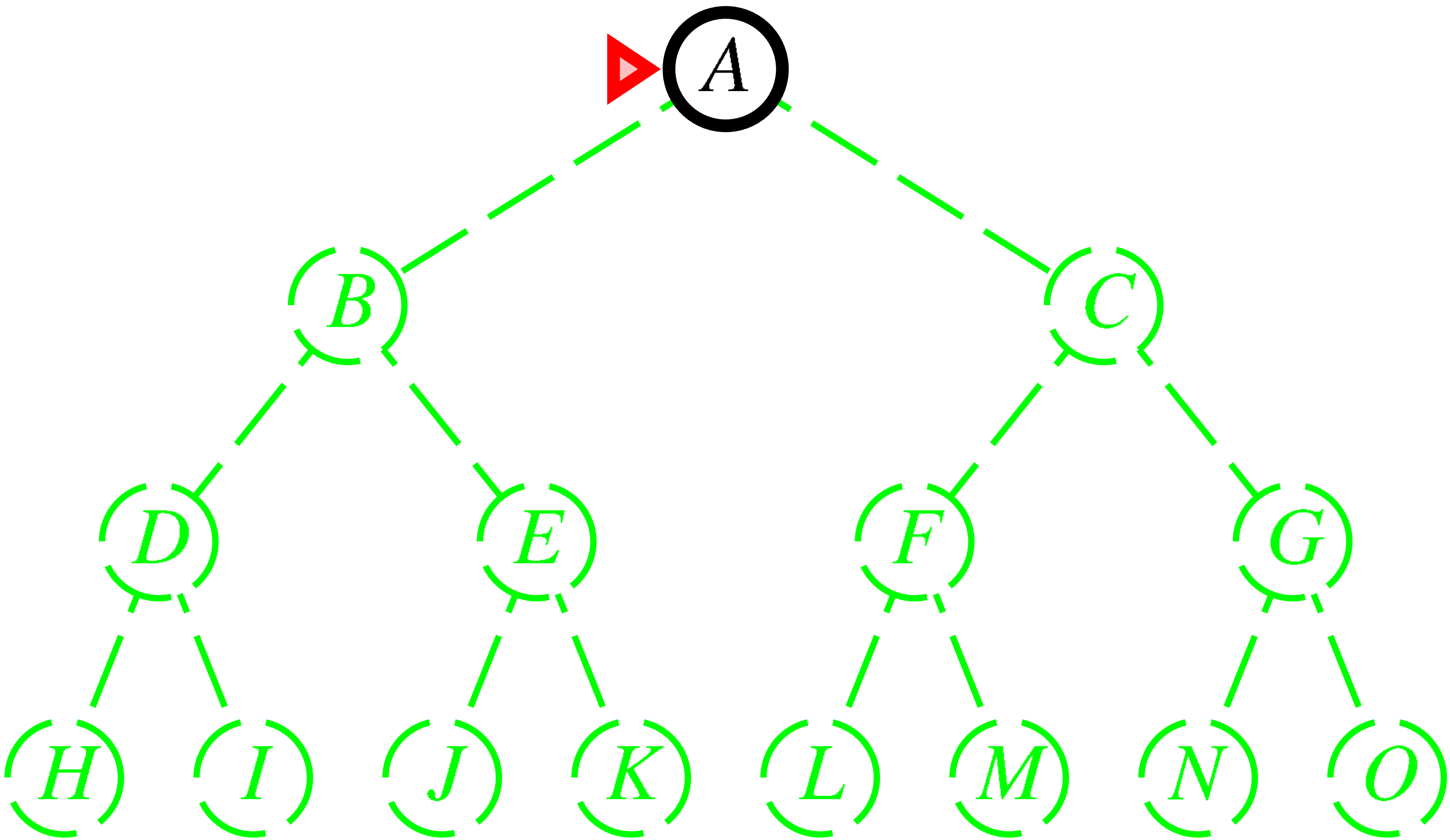
fringe **is a LIFO queue (a stack), i.e. successors go in at front of queue**

Note

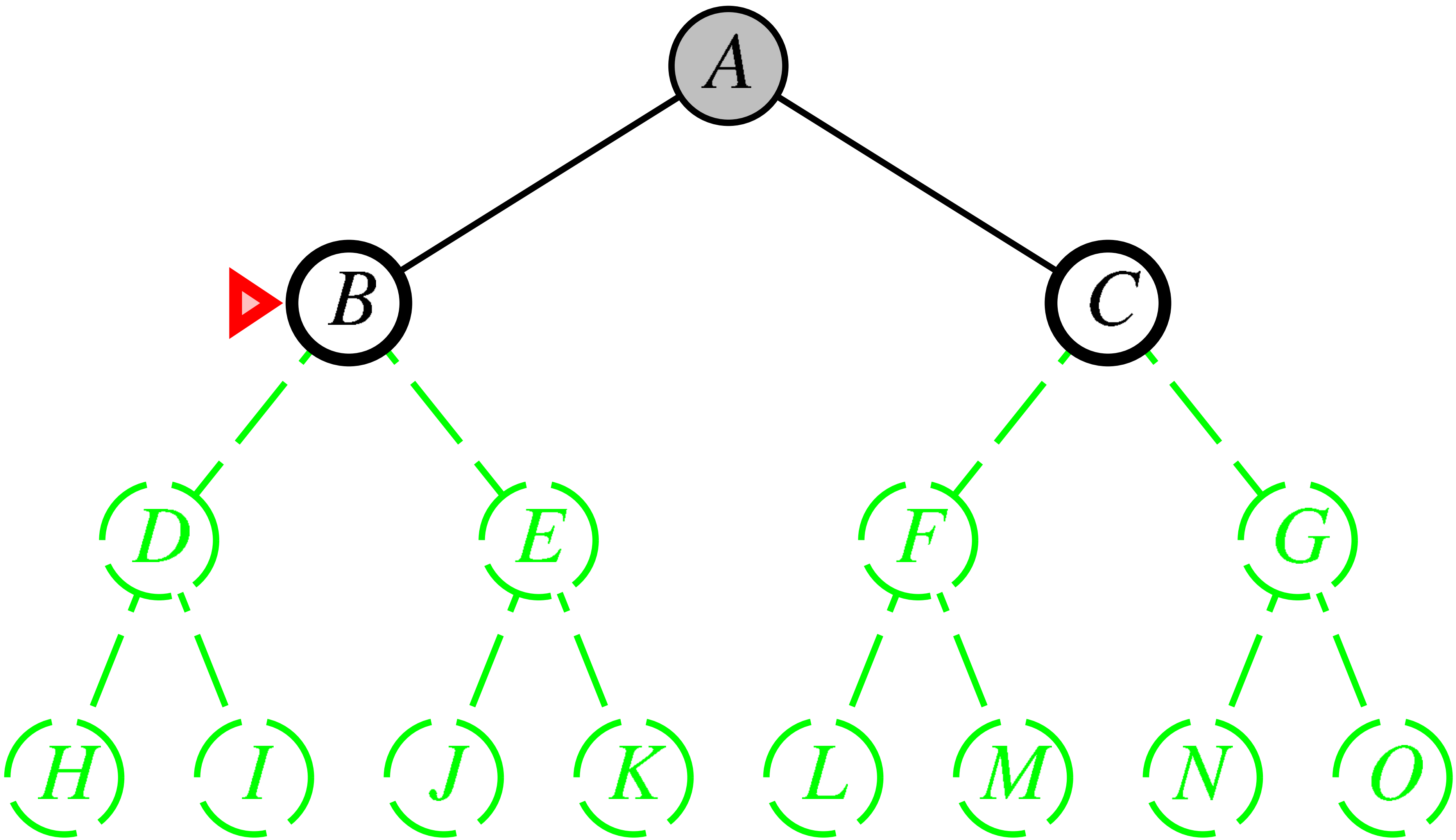
Depth-first search can perform infinite cyclic excursions

Need a finite, non-cyclic search space (or repeated-state checking)

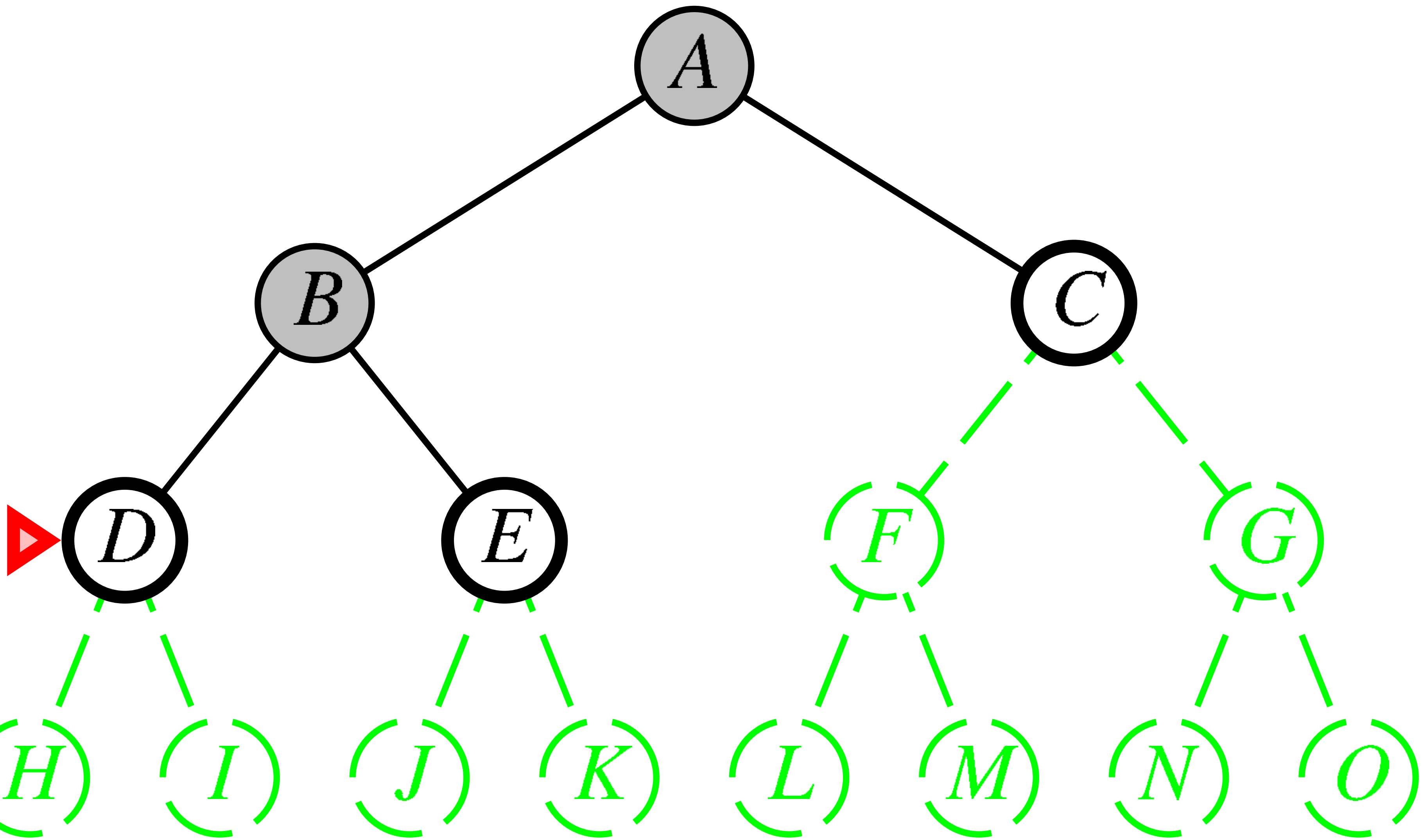
Depth-first search



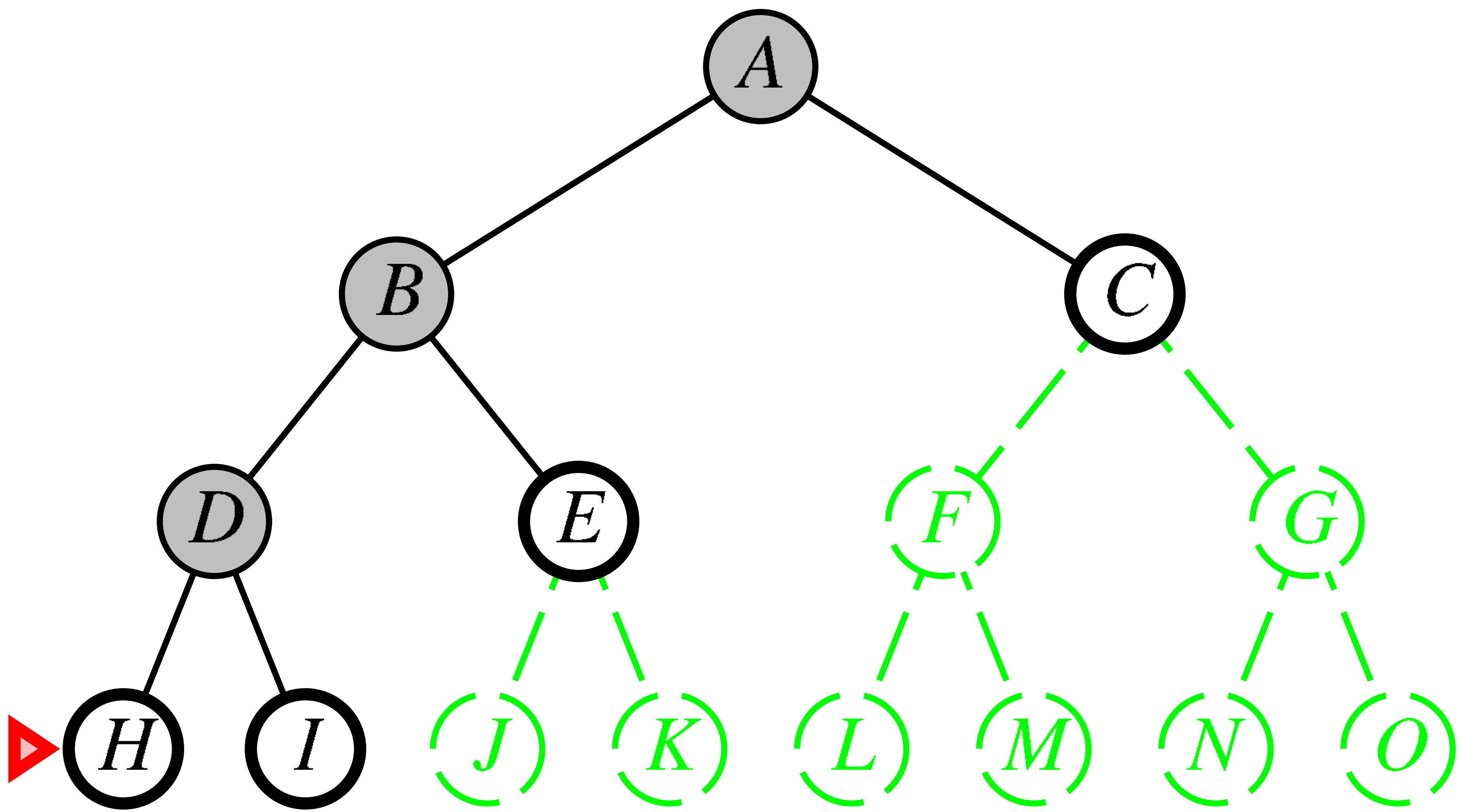
Depth-first search



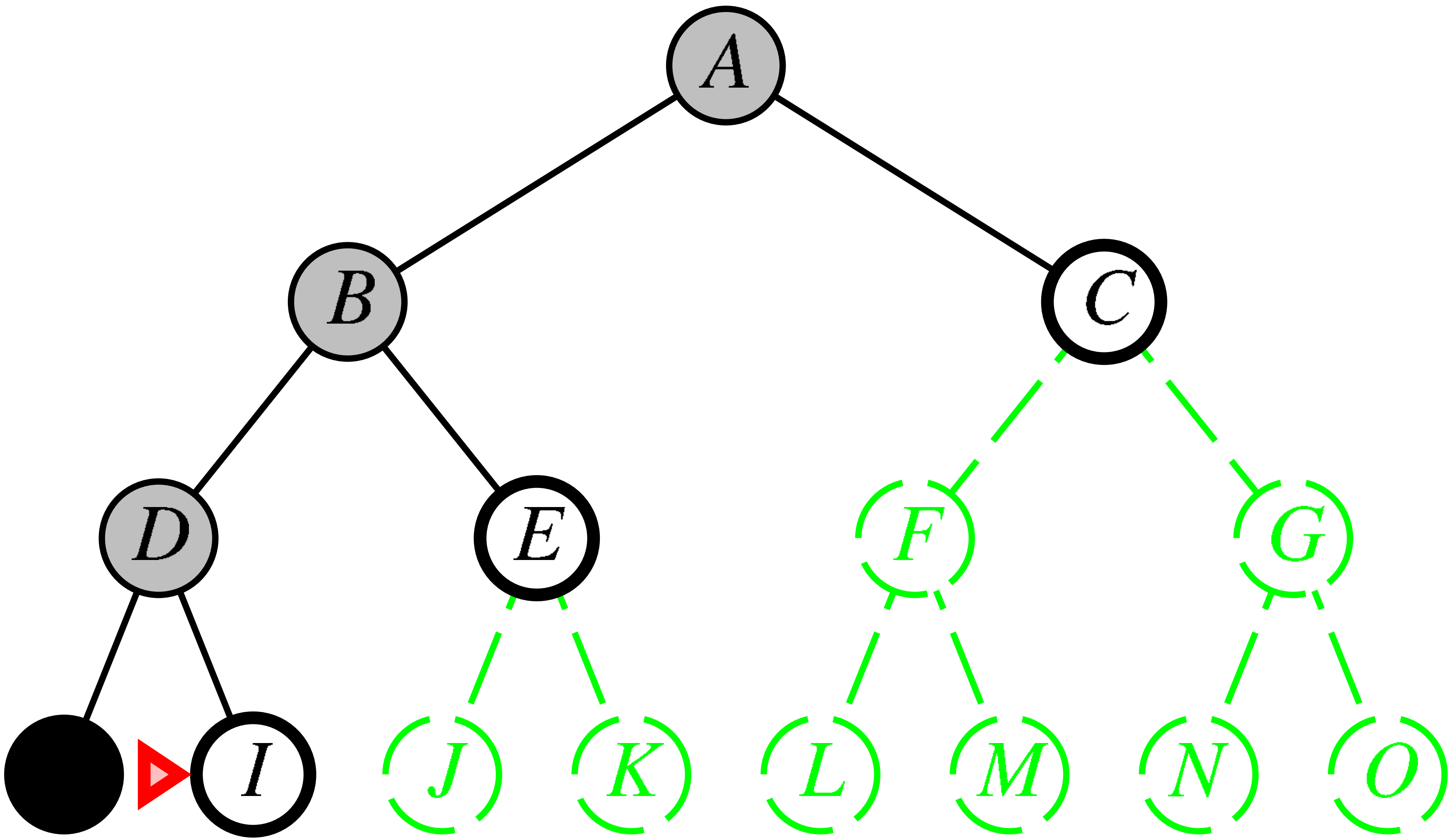
Depth-first search



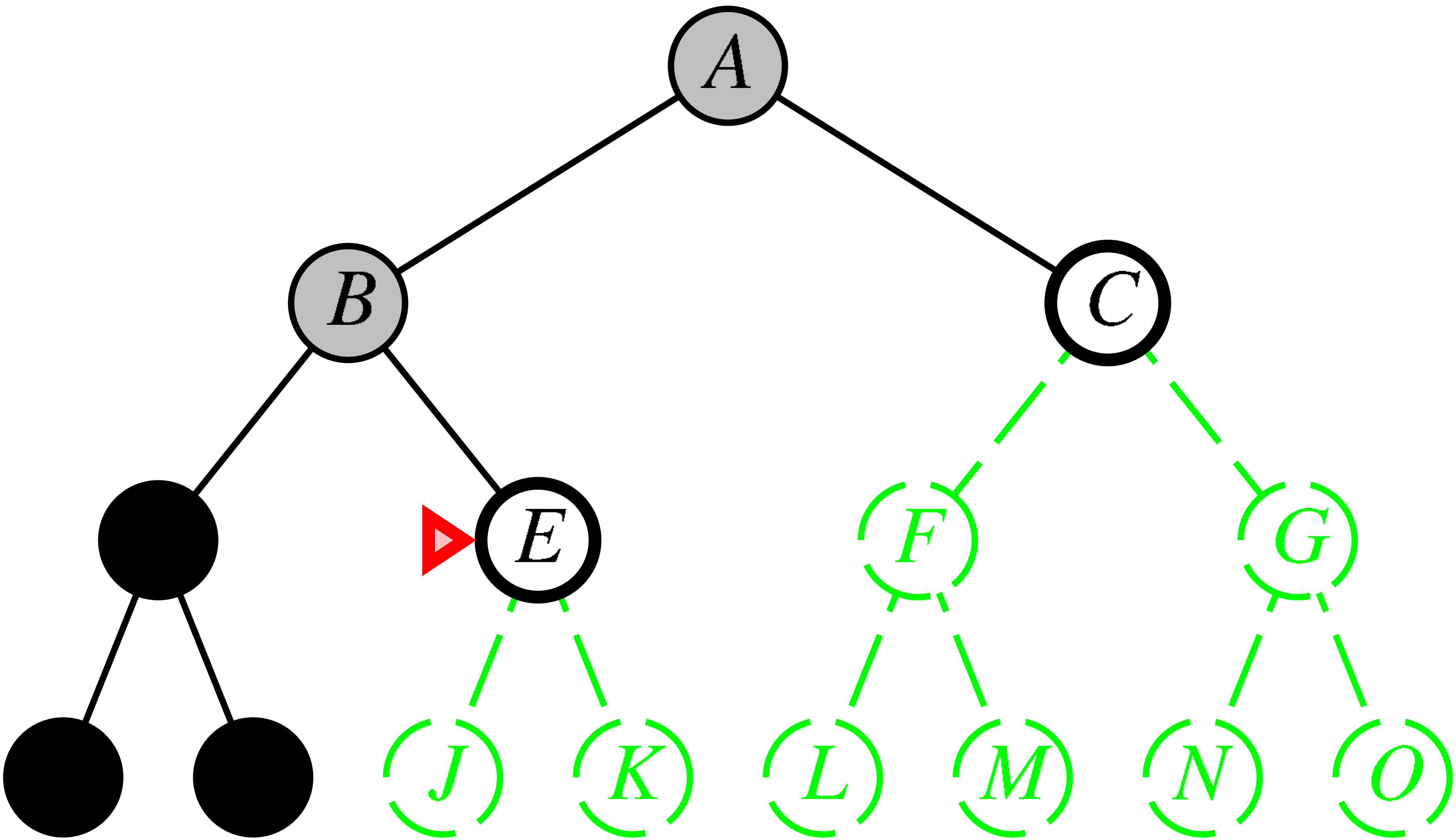
Depth-first search



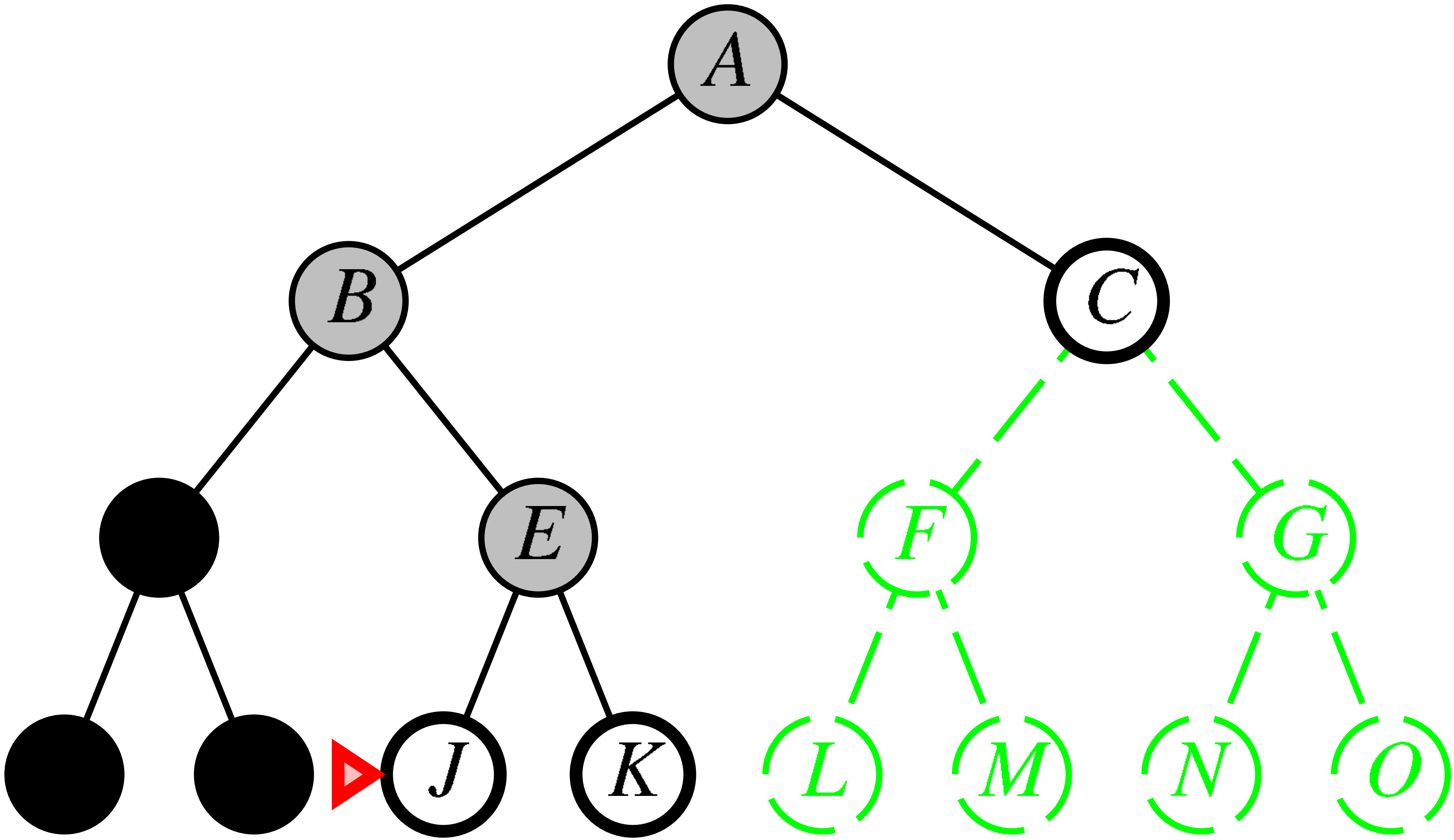
Depth-first search



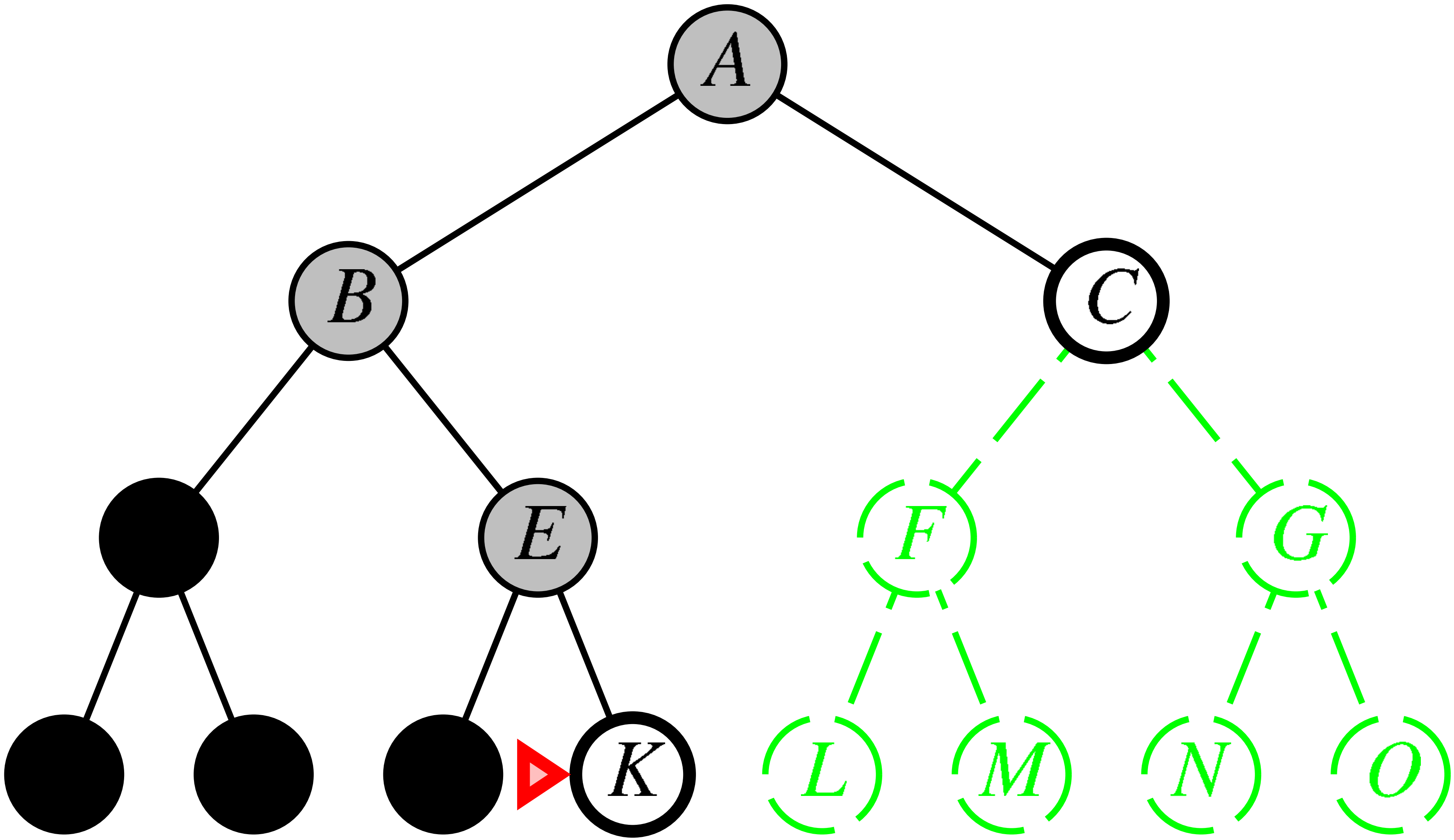
Depth-first search



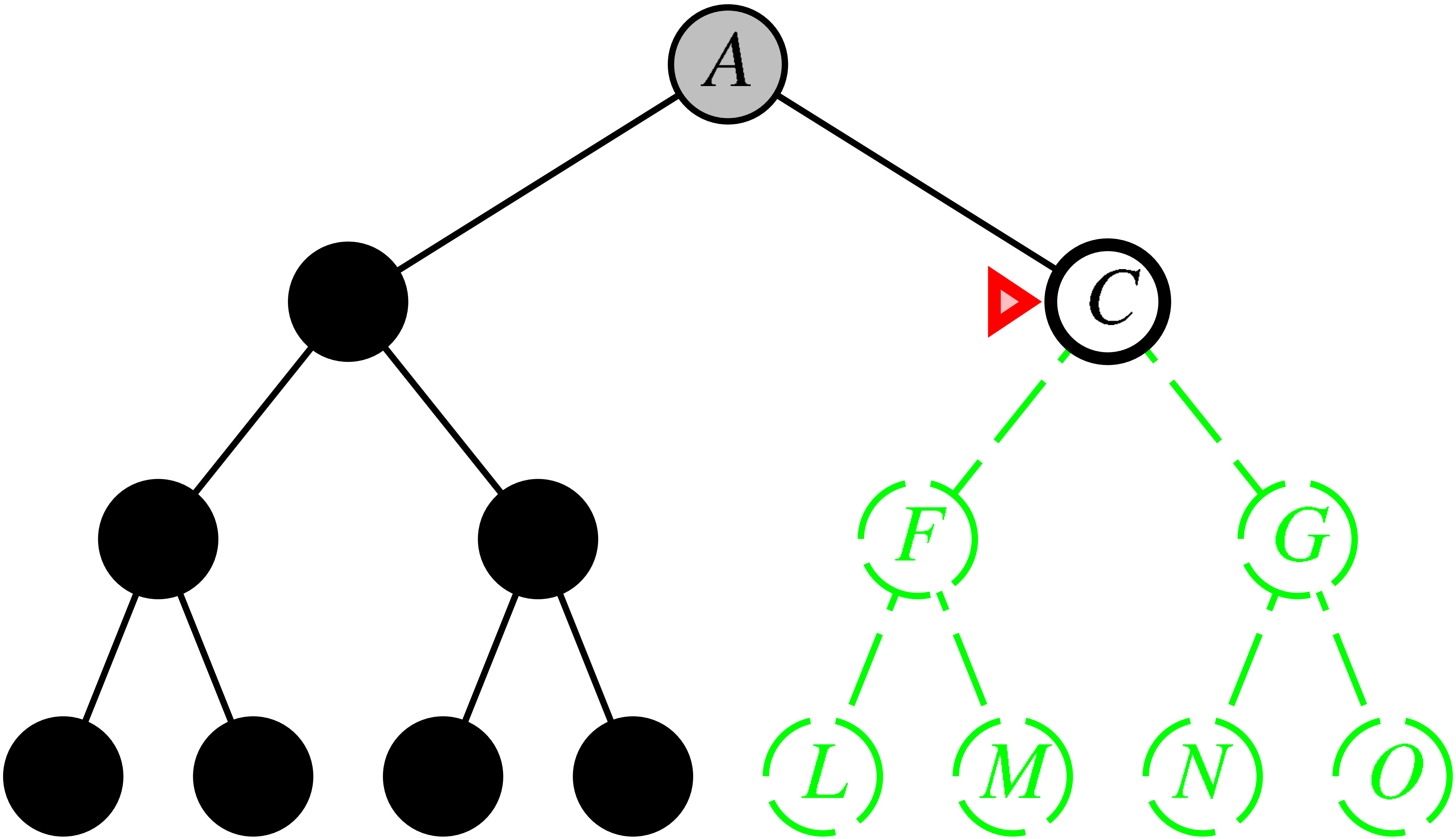
Depth-first search



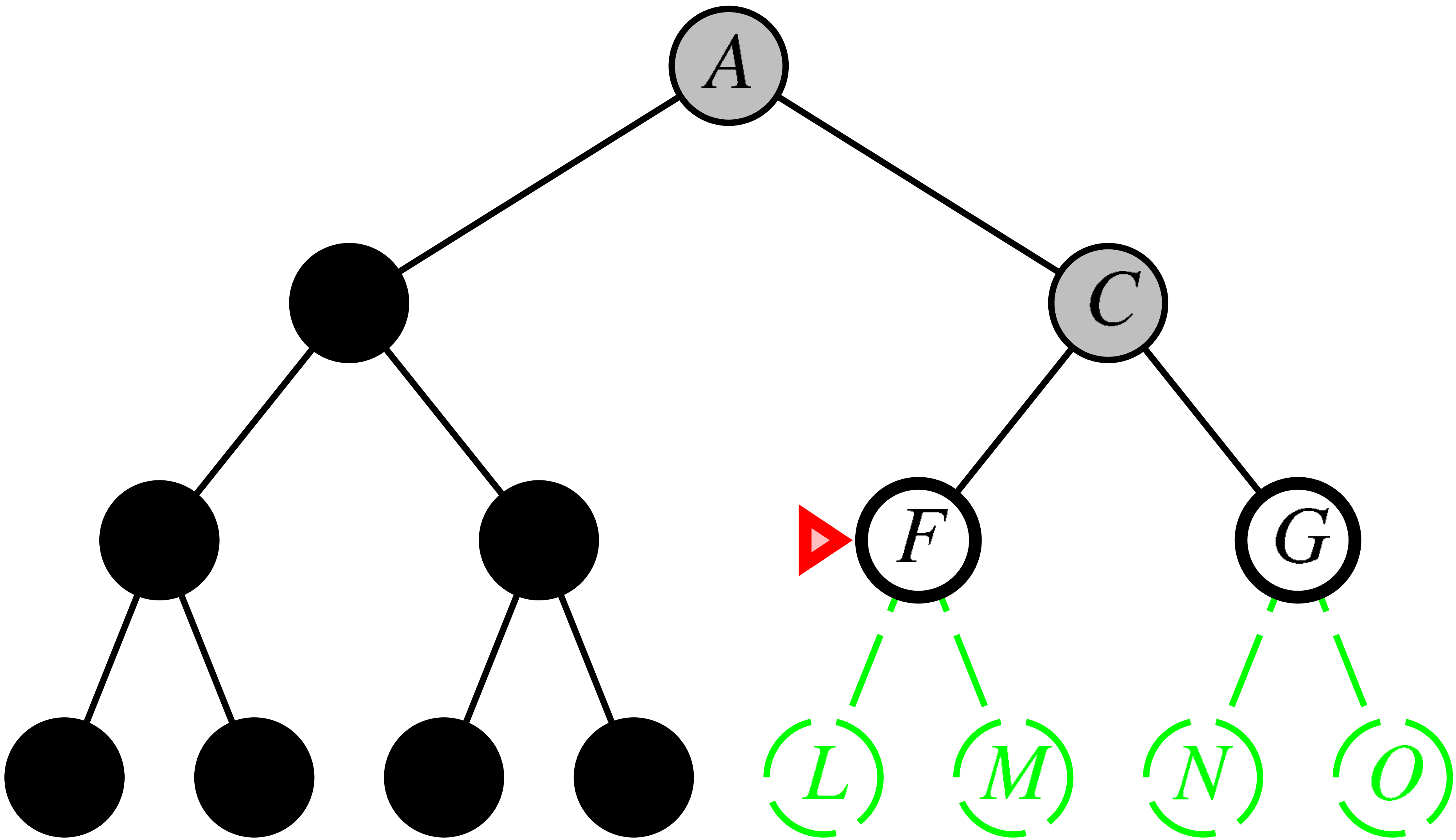
Depth-first search



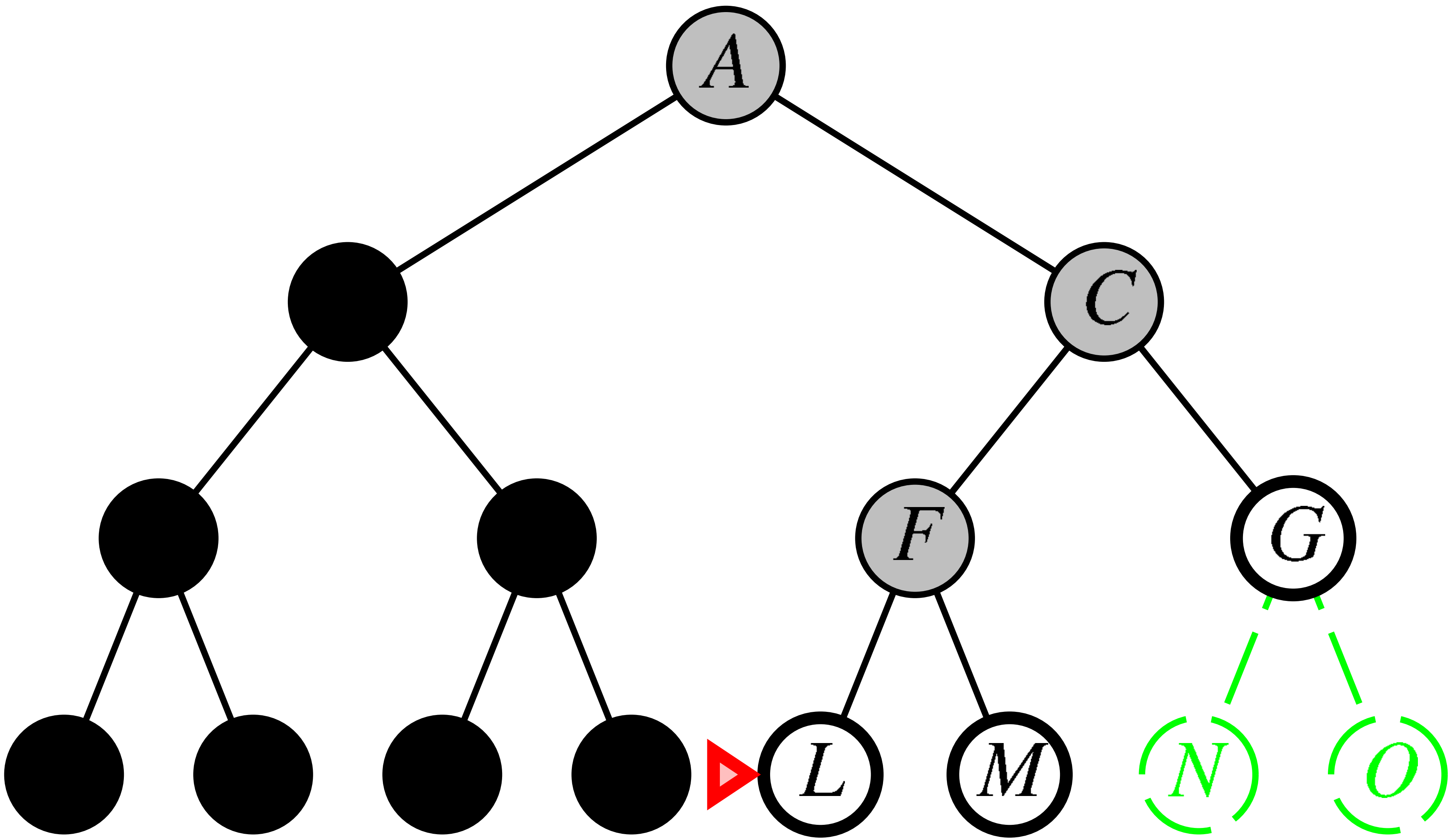
Depth-first search



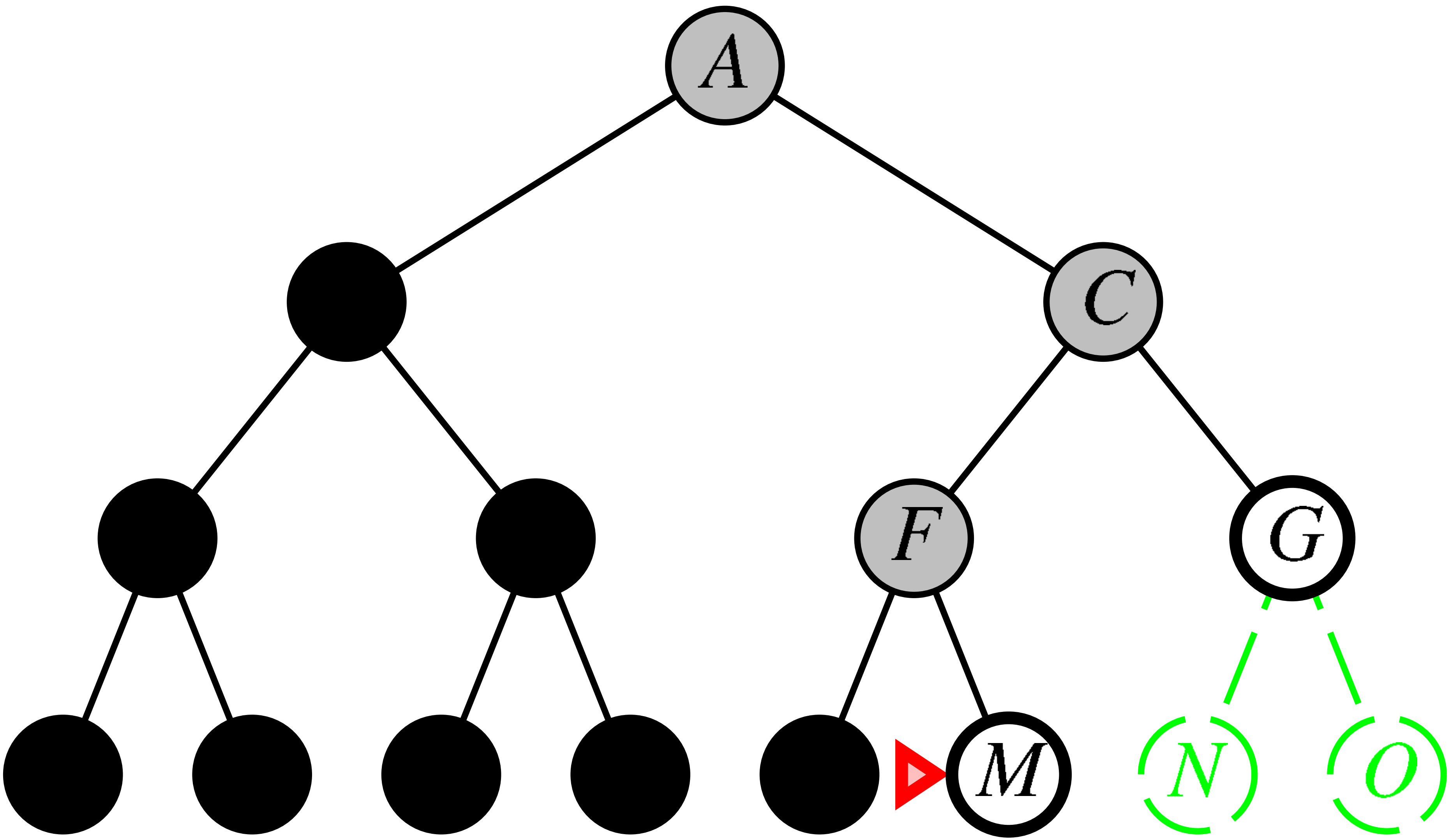
Depth-first search



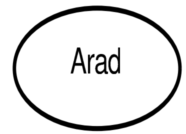
Depth-first search



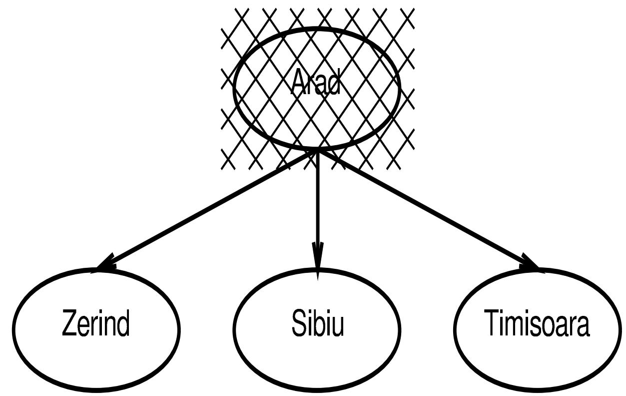
Depth-first search



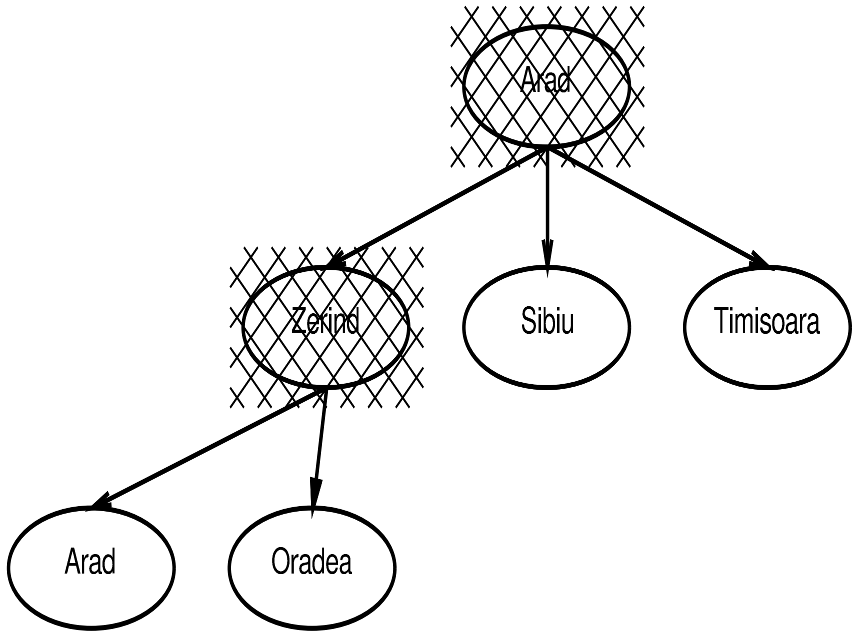
Depth-first search: Example Romania



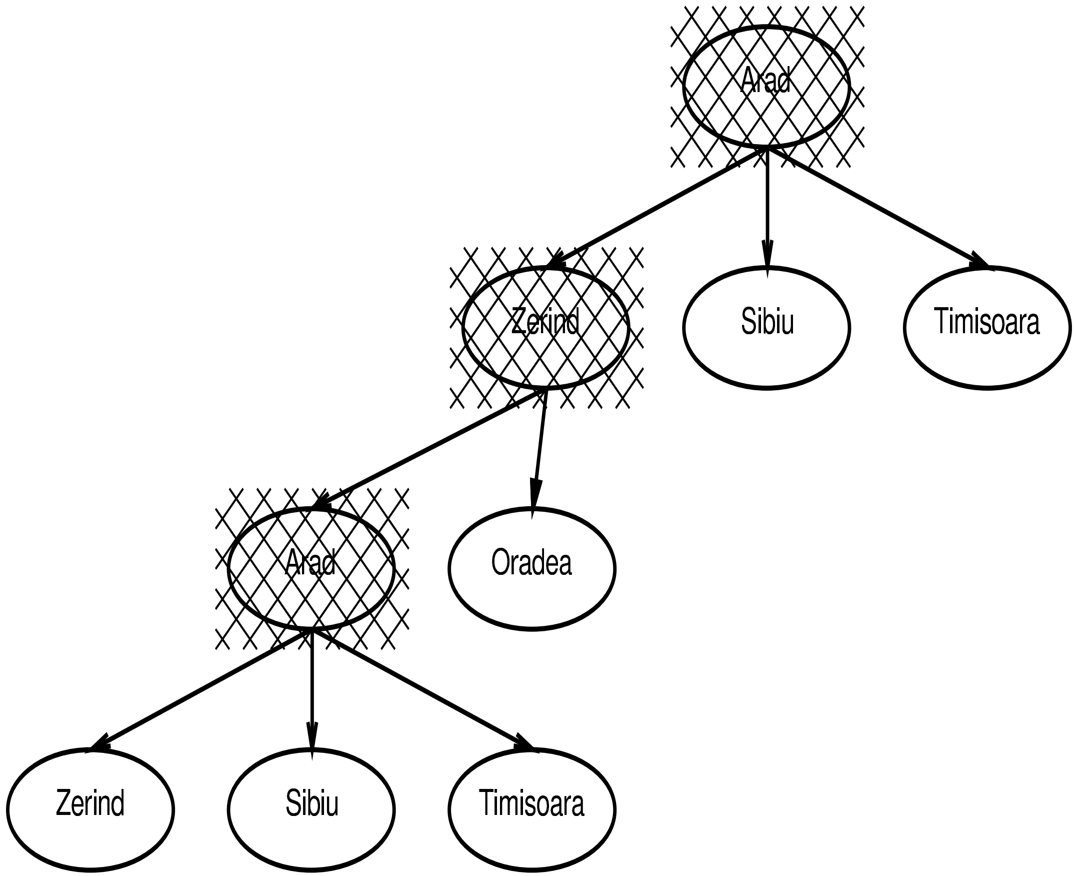
Depth-first search: Example Romania



Depth-first search: Example Romania



Depth-first search: Example Romania



Depth-first search: Properties

Complete

Time

Space

Optimal

Depth-first search: Properties

Complete Yes: if state space finite

No: if state contains infinite paths or loops

Time

Space

Optimal

Depth-first search: Properties

Complete Yes: if state space finite

No: if state contains infinite paths or loops

Time $O(b^m)$

Space

Optimal

Depth-first search: Properties

Complete Yes: if state space finite

No: if state contains infinite paths or loops

Time $O(b^m)$

Space $O(bm)$ (i.e. linear space)

Optimal

Depth-first search: Properties

Complete	**Yes:** if state space finite
	No: if state contains infinite paths or loops
Time	$O(b^m)$
Space	$O(bm)$ **(i.e. linear space)**
Optimal	**No**

Depth-first search: Properties

Complete — **Yes:** if state space finite

No: if state contains infinite paths or loops

Time — $O(b^m)$

Space — $O(bm)$ (i.e. linear space)

Optimal — No

Disadvantage

Time terrible if m much larger than d

Advantage

Time may be much less than breadth-first search if solutions are dense

Iterative deepening search

Depth-limited search

Depth-first search with depth limit

Iterative deepening search

Depth-limited search

Depth-first search with depth limit

Iterative deepening search

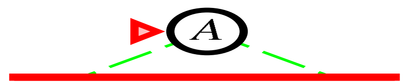
Depth-limit search with ever increasing limits

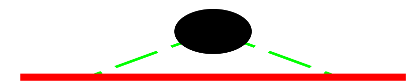
function ITERATIVE-DEEPENING-SEARCH(*problem*) **returns** a solution or failure
    inputs: *problem*   /* a problem */

    for *depth* $\leftarrow$ 0 **to** ∞ **do**
        result $\leftarrow$ DEPTH-LIMITED-SEARCH(*problem*, *depth*)
        if *result* $\neq$ cutoff **then return** *result*
    end

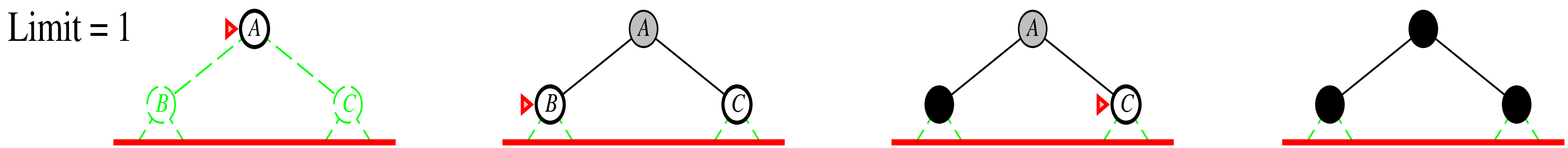
Iterative deepening search with depth limit 0

$\text{Limit} = 0$

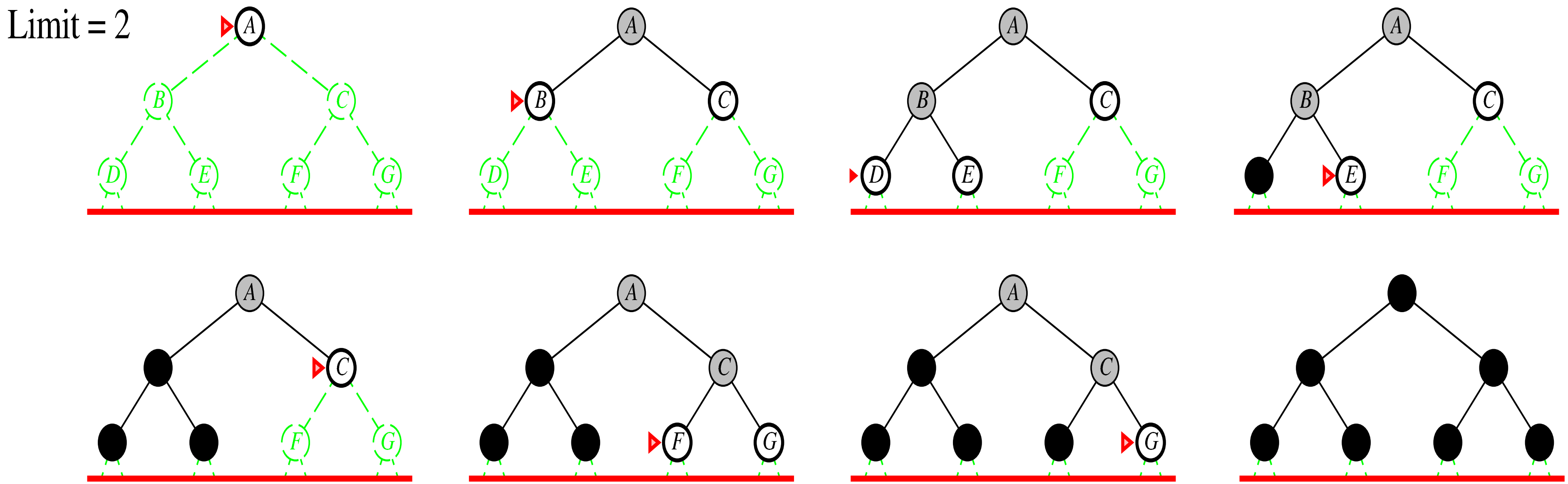




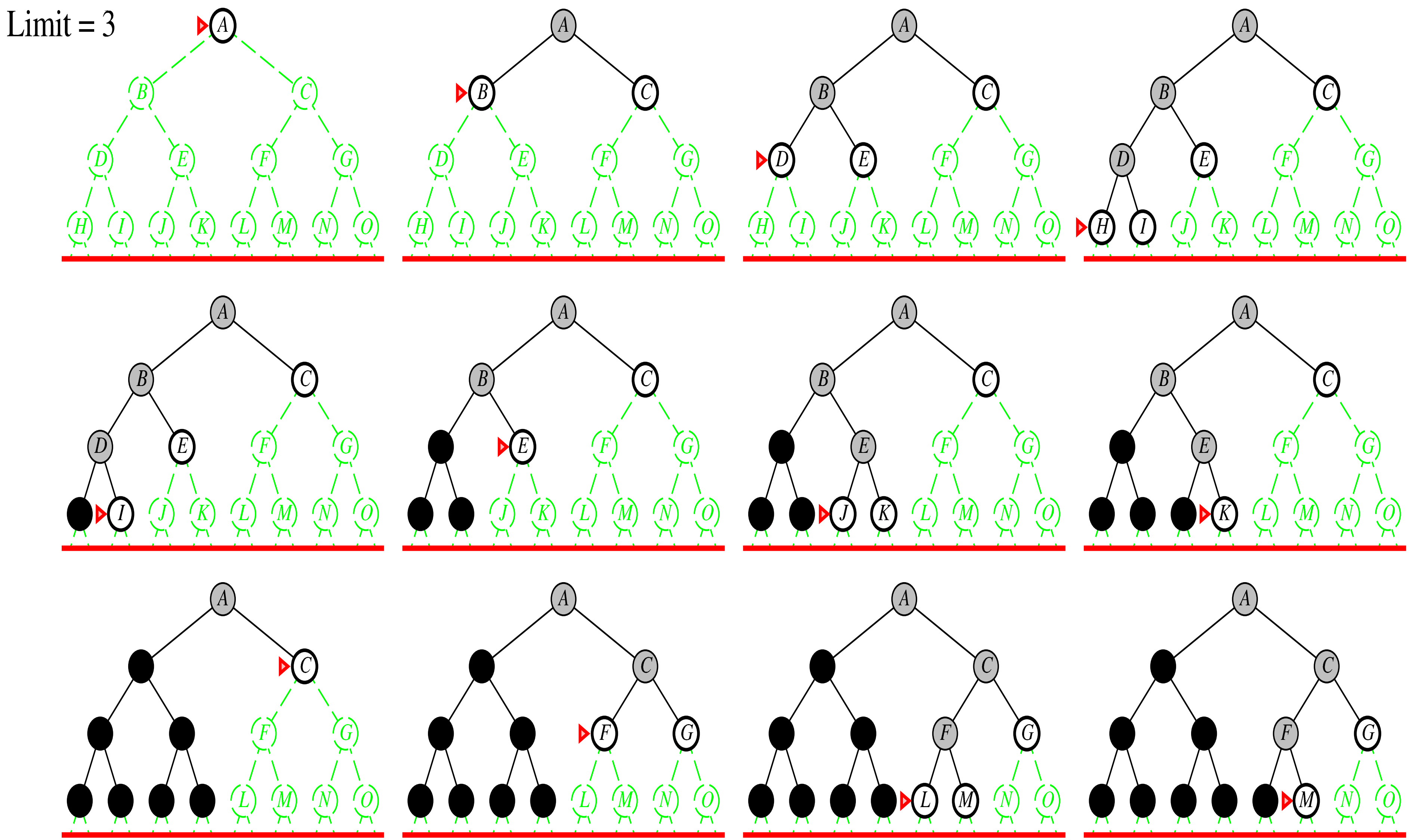
Iterative deepening search with depth limit 1



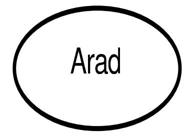
Iterative deepening search with depth limit 2



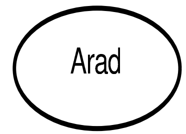
Iterative deepening search with depth limit 3



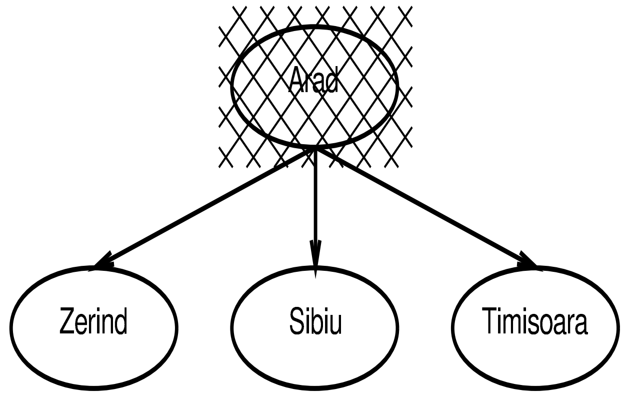
Iterative deepening search: Example Romania with $l = 0$



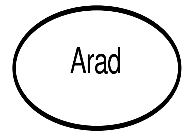
Iterative deepening search: Example Romania with $l = 1$



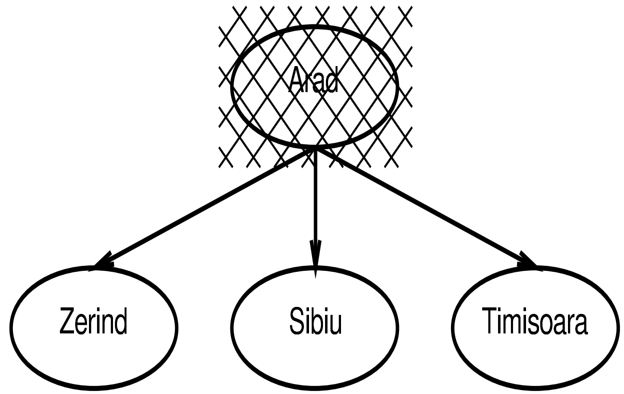
Iterative deepening search: Example Romania with $l = 1$



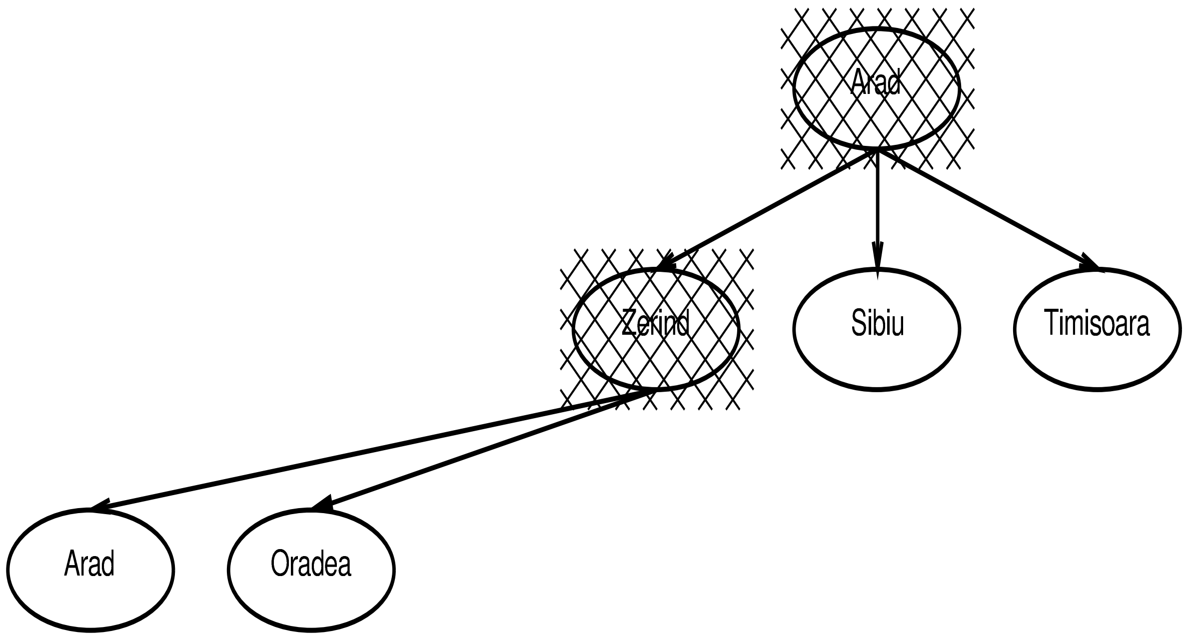
Iterative deepening search: Example Romania with $l = 2$



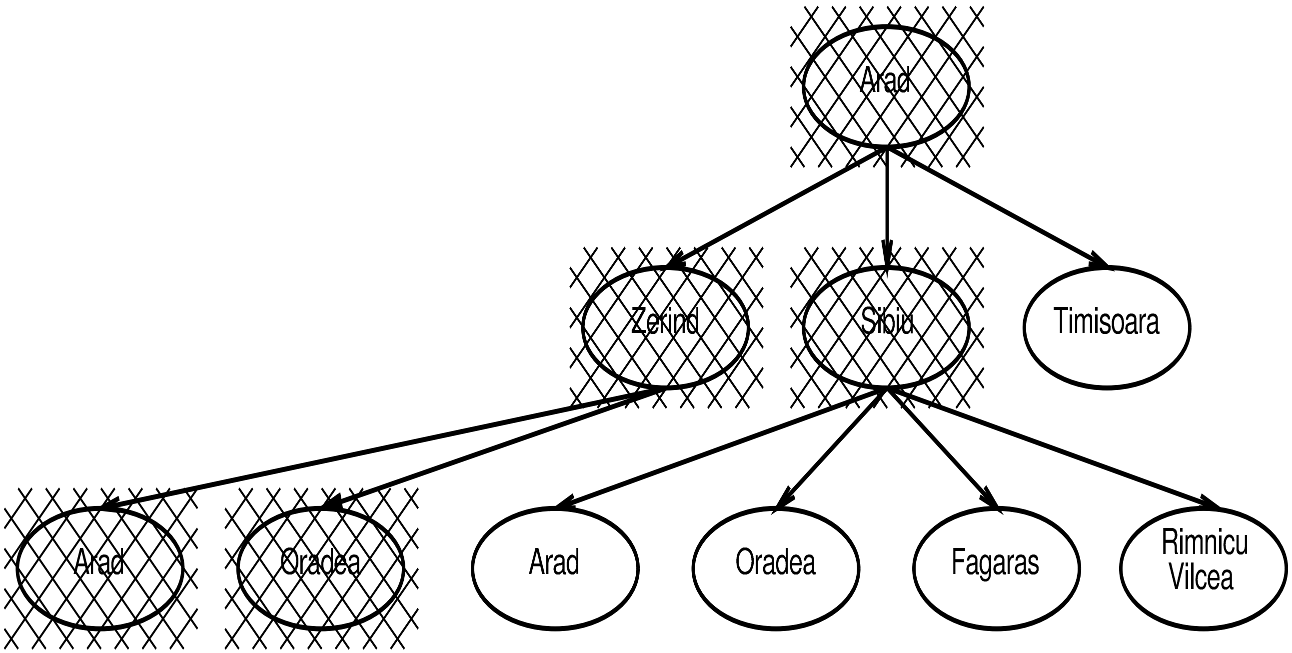
Iterative deepening search: Example Romania with $l = 2$



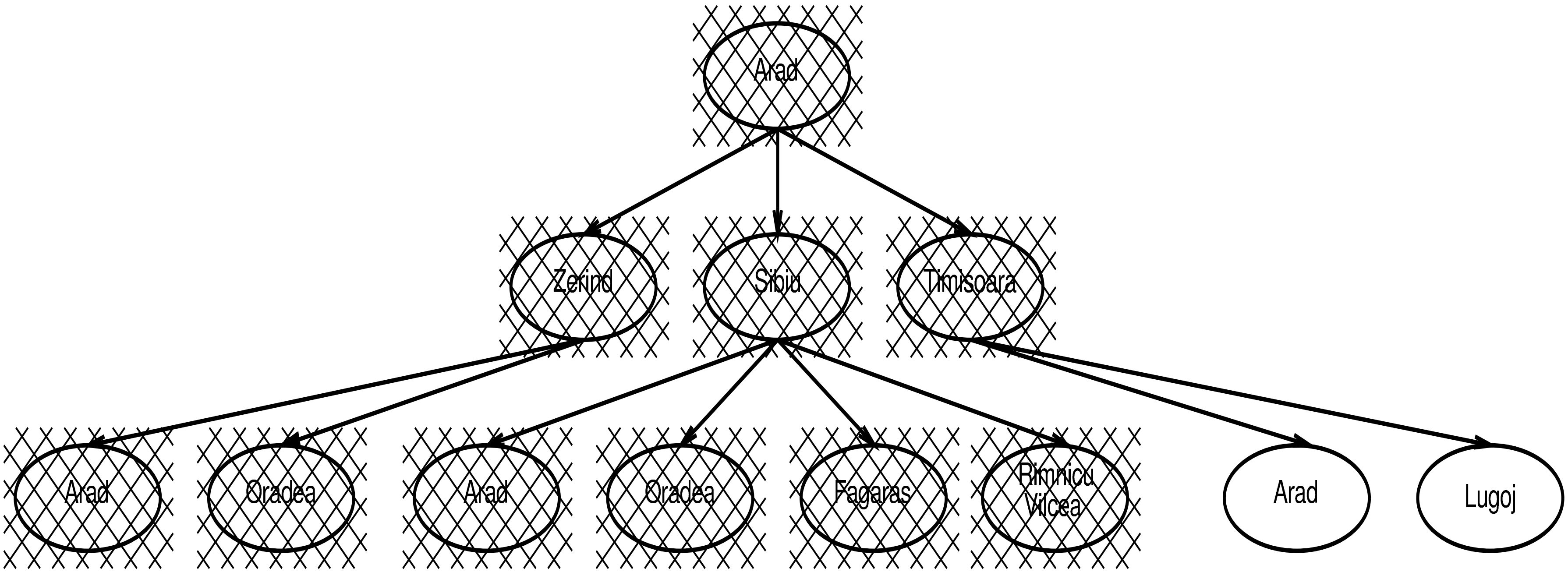
Iterative deepening search: Example Romania with $l = 2$



Iterative deepening search: Example Romania with $l = 2$



Iterative deepening search: Example Romania with $l = 2$



Iterative deepening search: Properties

Complete

Time

Space

Optimal

Iterative deepening search: Properties

Complete Yes

Time

Space

Optimal

Iterative deepening search: Properties

Complete Yes

Time $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d \in O(b^{d+1})$

Space

Optimal

Iterative deepening search: Properties

Complete Yes

Time $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d \in O(b^{d+1})$

Space $O(bd)$

Optimal

Iterative deepening search: Properties

Complete **Yes**

Time $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d \in O(b^{d+1})$

Space $O(bd)$

Optimal **Yes (if step cost = 1)**

Iterative deepening search: Properties

Complete	**Yes**
Time	$(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d \in O(b^{d+1})$
Space	$O(bd)$
Optimal	**Yes (if step cost = 1)**

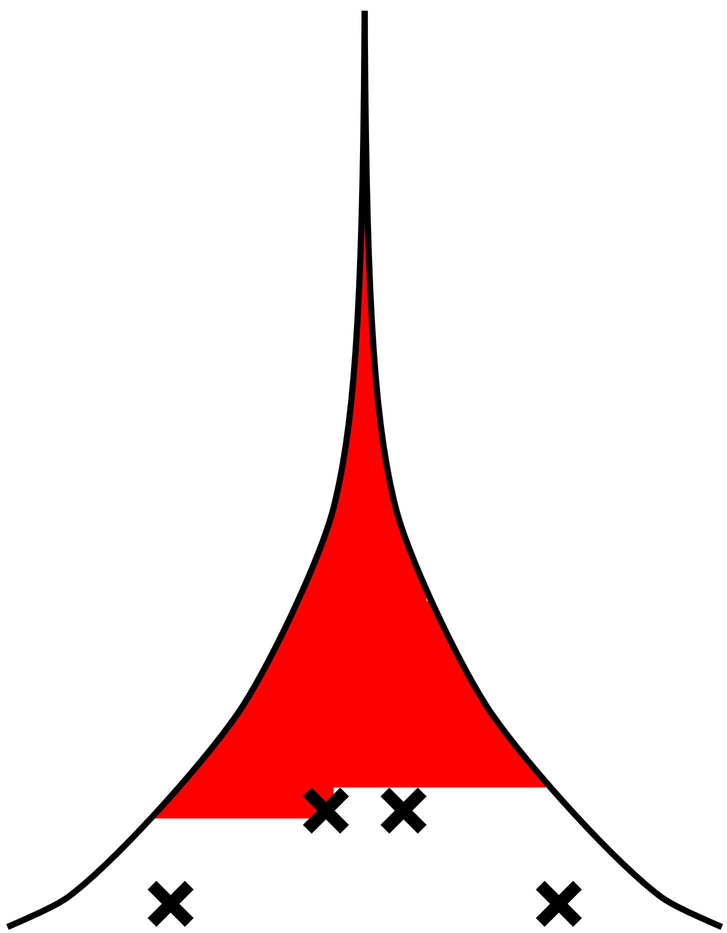
(Depth-First) Iterative-Deepening Search often used in practice for search spaces of large, infinite, or unknown depth.

Comparison

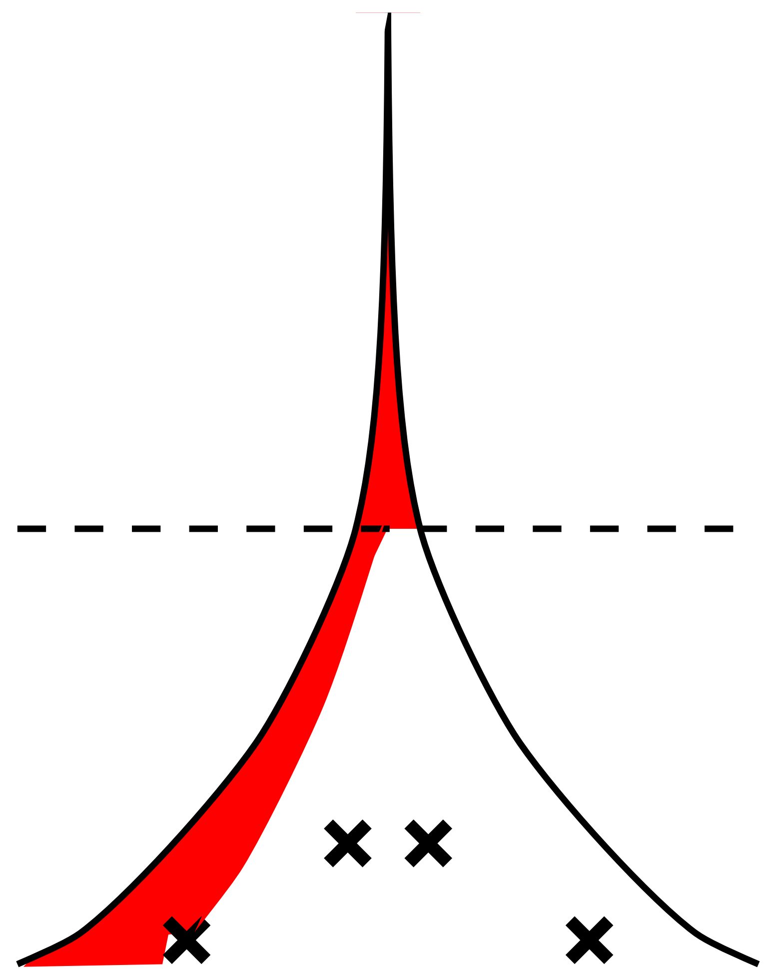
Criterion	Breadth-first	Uniform-cost	Depth-first	Iterative deepening
Complete?	**Yes***	**Yes***	**No**	**Yes**
Time	b^{d+1}	$\approx b^d$	b^m	b^d
Space	b^{d+1}	$\approx b^d$	bm	bd
Optimal?	**Yes***	**Yes**	**No**	**Yes**

Comparison

Breadth-first search



Iterative deepening search



Summary

- Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored
- Variety of uninformed search strategies
- Iterative deepening search uses only linear space and not much more time than other uninformed algorithms